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- 35 institutes, by the *Shanghai Jiao Tong University (China)* in 2008
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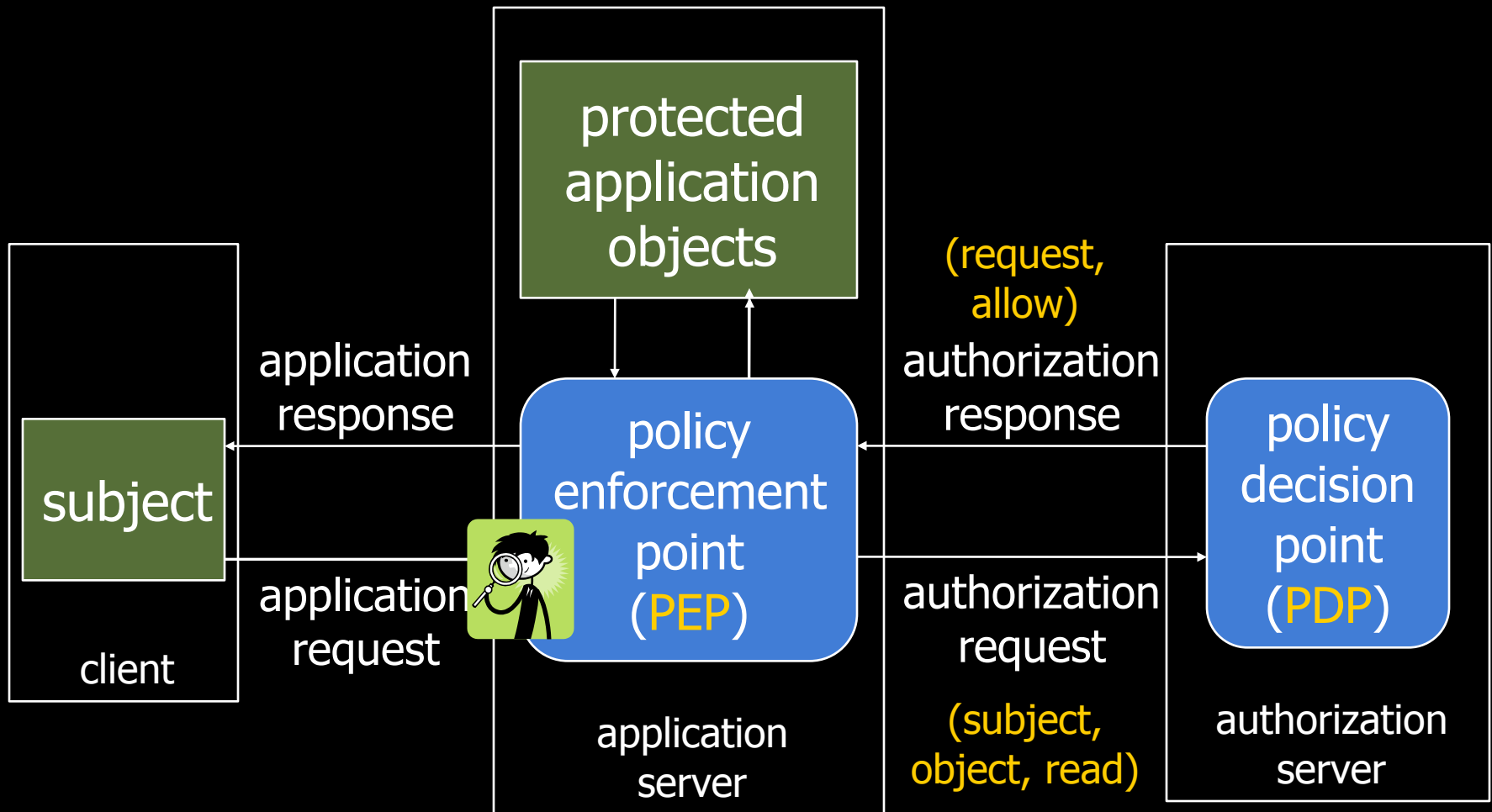


Toward Improving Availability and Performance of Enterprise Authorization Services

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typical authorization architecture



IBM Access Manager, Entrust GetAccess, CA SiteMinder, etc.

request-response model

request-response model



- + re-use of authorization logic
- + consistent policy enforcement
- + lower admin overhead



- reduced availability
- increased latency
- reduced scalability

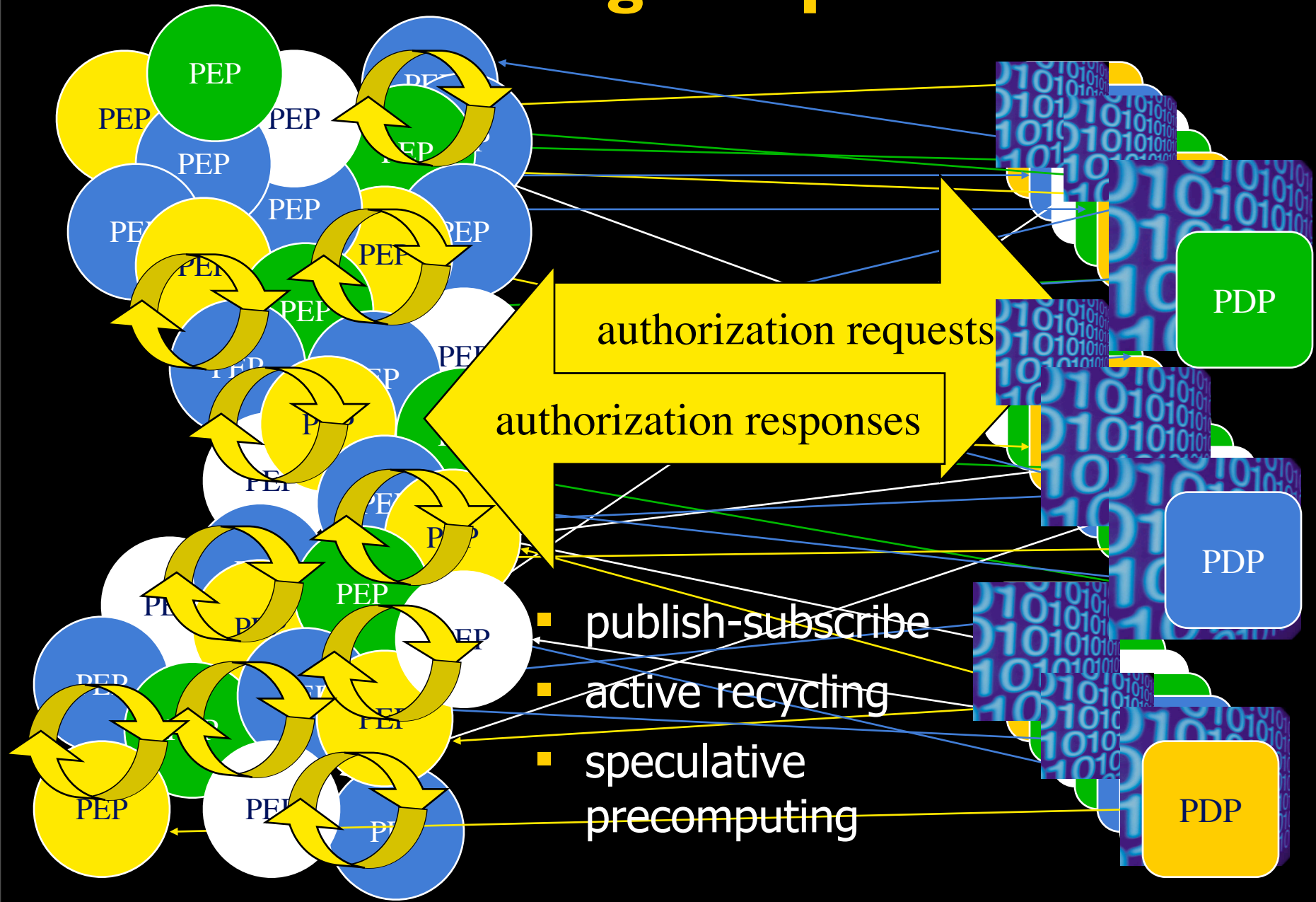
departing assumptions

1. **processor** resources virtually **free**
2. **commodity** computing most **cost-effective**
3. network **bandwidth** virtually **unlimited**
4. **human** time/attention **expensive**

existing approaches

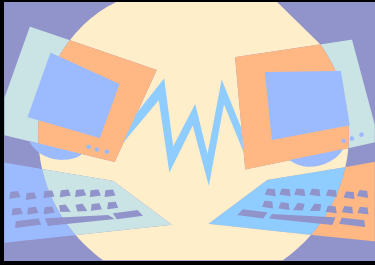
- caching previous authorizations
 - + simple, inexpensive
 - + improves performance & availability
 - serves only returning requests (precise recycling)
- generic fault-tolerance through replication/redundancy
 - + improve availability
 - latency remains unchanged
 - require specialized OS/middleware
 - poorly scale on large populations

addressing the problem



outline

- authorization architecture based on pub-sub
- concept and model for inferring new authorizations from previous responses:
secondary and approximate authorization model (SAAM)
- SAAM algorithms for BLP and RBAC
- distributed and cooperative SAAM

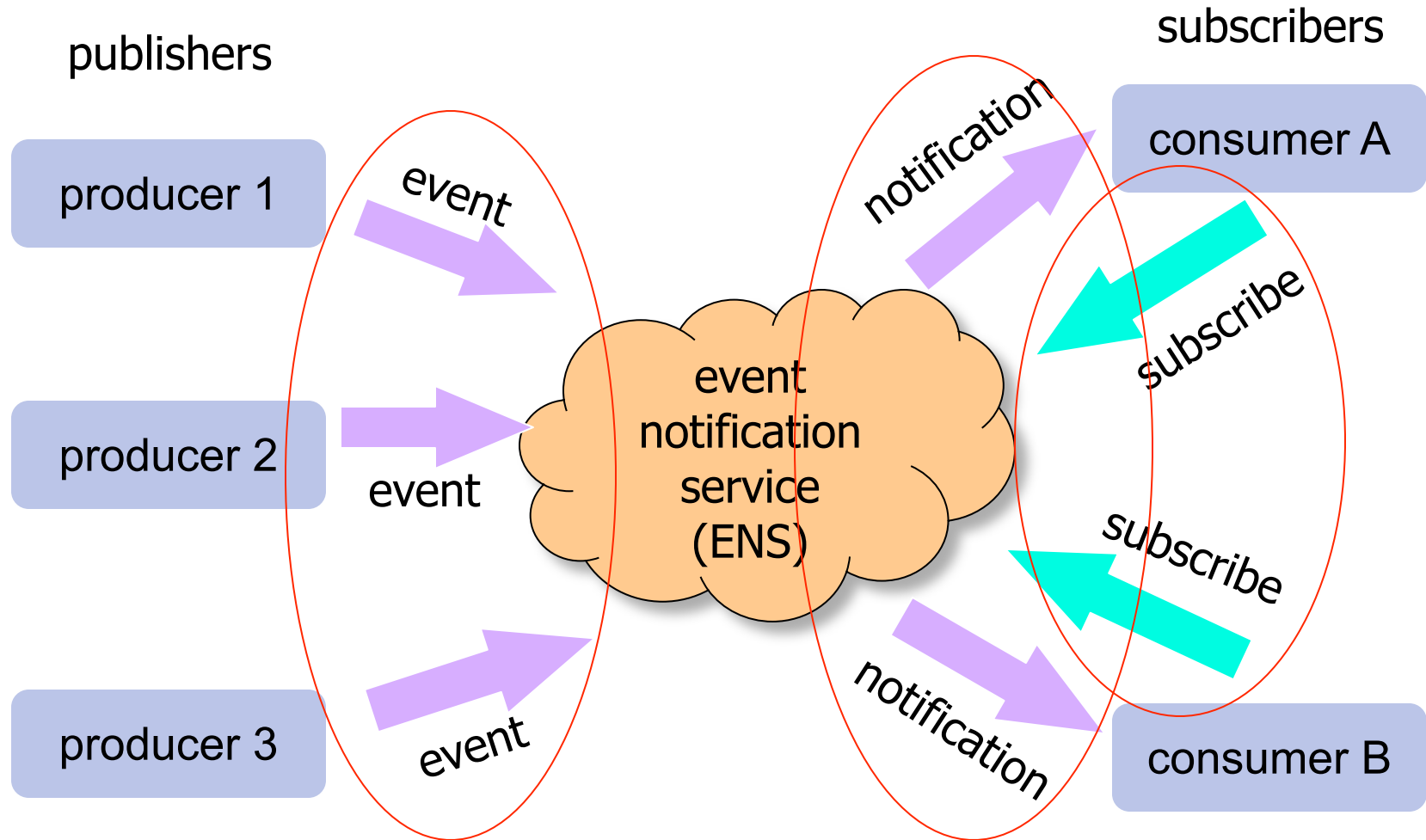


PUB-SUB

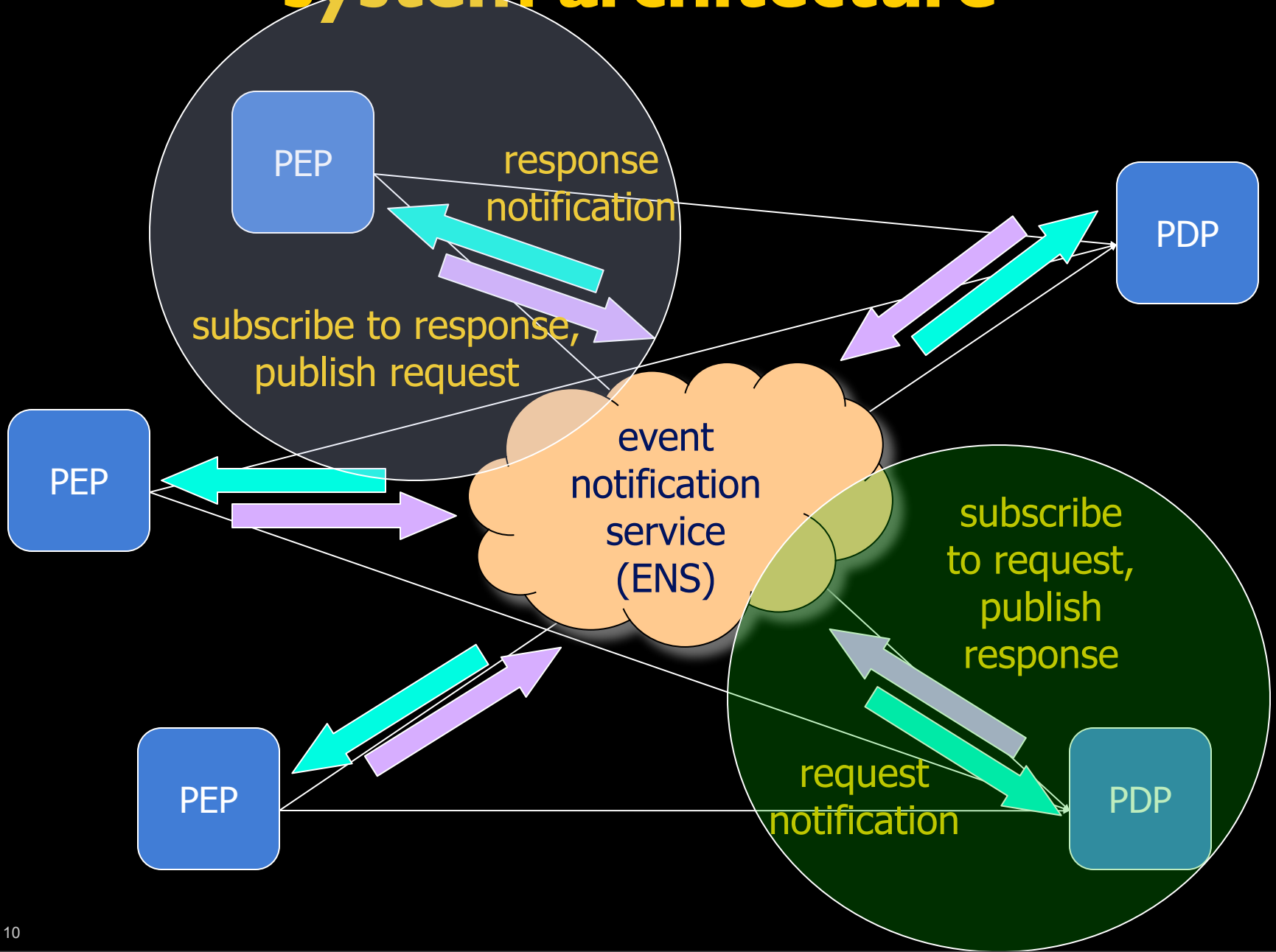
authorization architecture based on publish-subscribe model

Q. Wei, M. Ripeanu, K. Beznosov “**Authorization using Publish-Subscribe Model,**”
in *Proceedings of the IEEE International Symposium on Parallel and Distributed
Processing with Applications (ISPA'08)*, December 10-12, 2008, Sydney, Australia, pp.
53-62

basic components in a publish/subscribe system

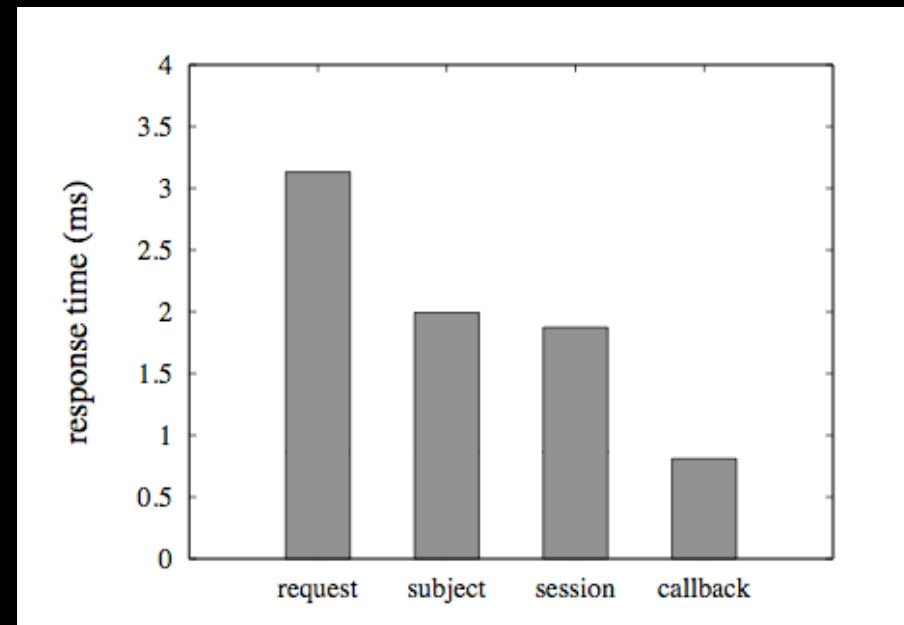
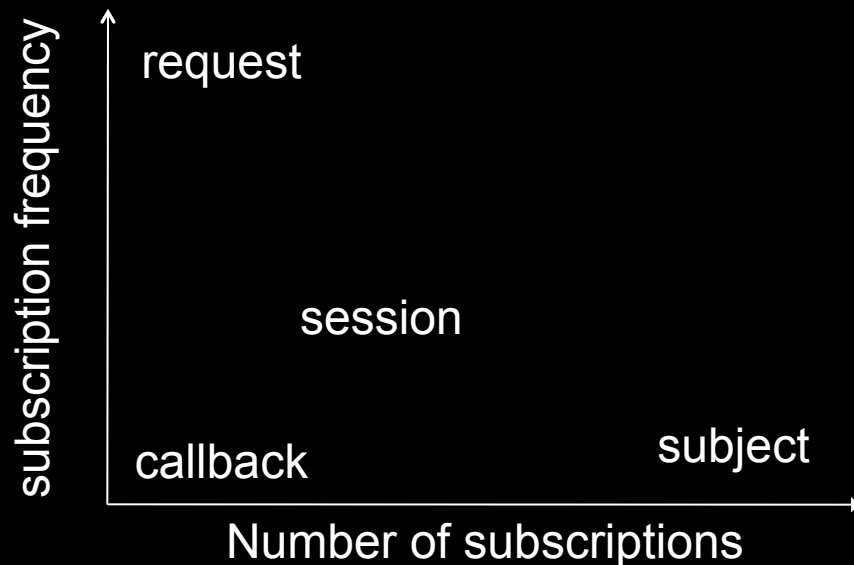


system architecture



PEP subscription schemes

- per request
- per subject
- per user session
- decisions delivered via callback, instead of pub-sub



PDP subscription schemes

- makes all the subscriptions at start-up
 - subscription frequency is zero
- can subscribe to
 - all requests
 - all resources
 - resource groups
 - application groups
- other options
 - subject groups

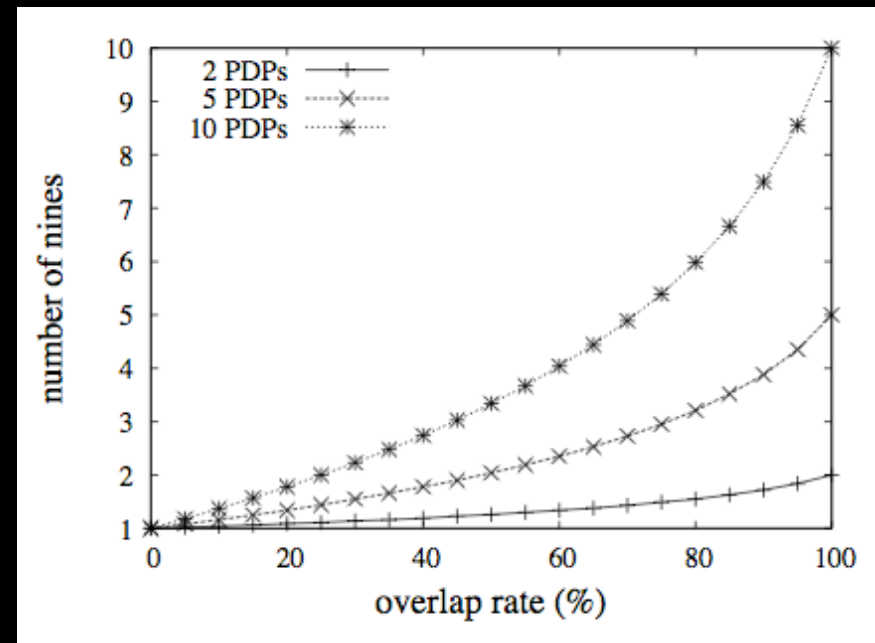
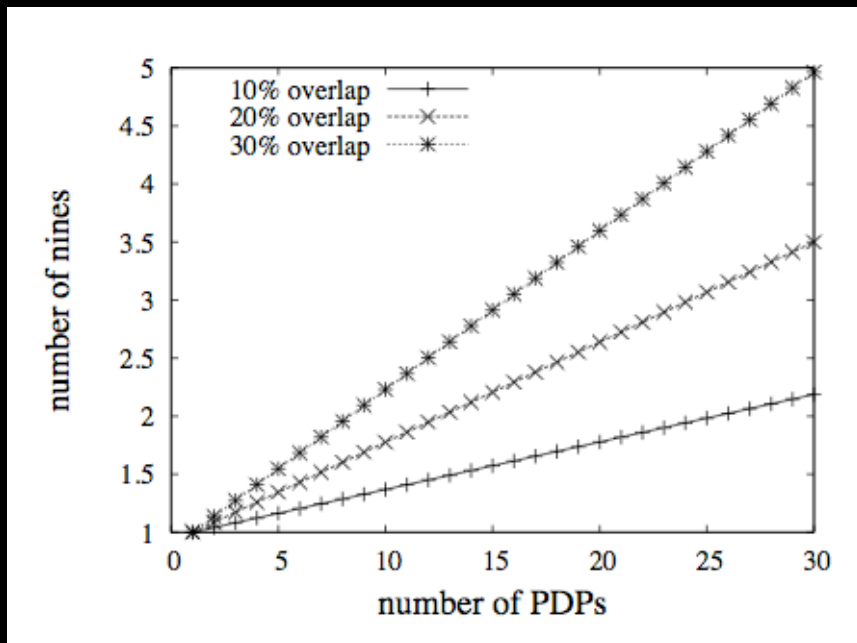
↑
number of subscriptions

evaluation

availability analysis

$$p_t = 1 - (1 - p) (1 - p \cdot o)^{m-1}$$

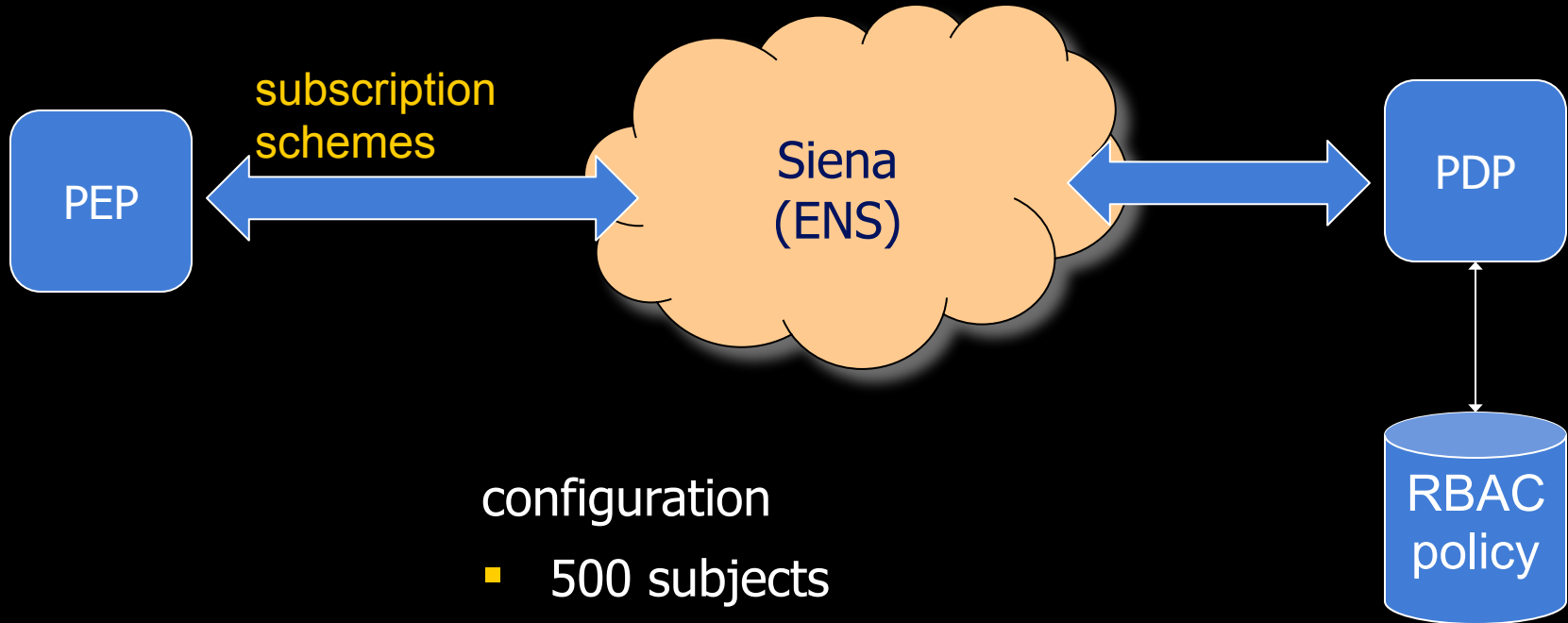
- p - availability of each PDP
- m - number of PDPs
- o - overlap of served request spaces



performance evaluation

- metrics
 - response time
 - maximum throughput
- influencing factors
 - number of subscriptions
 - subscription frequency

prototype

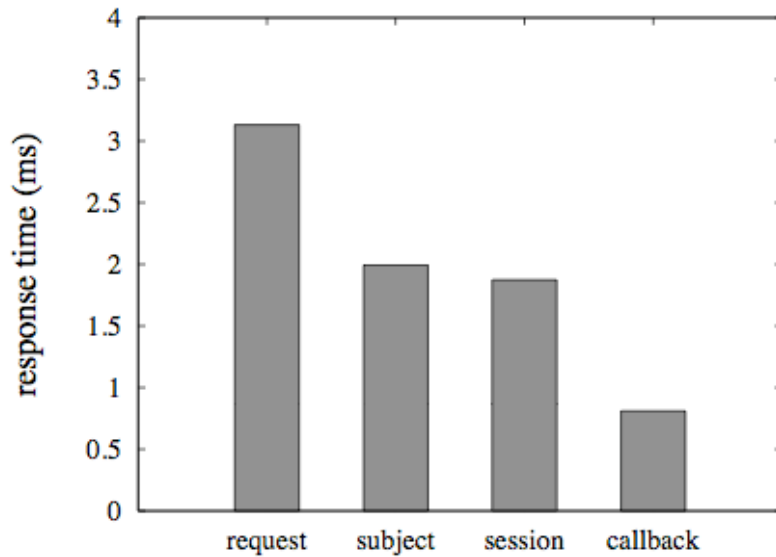


configuration

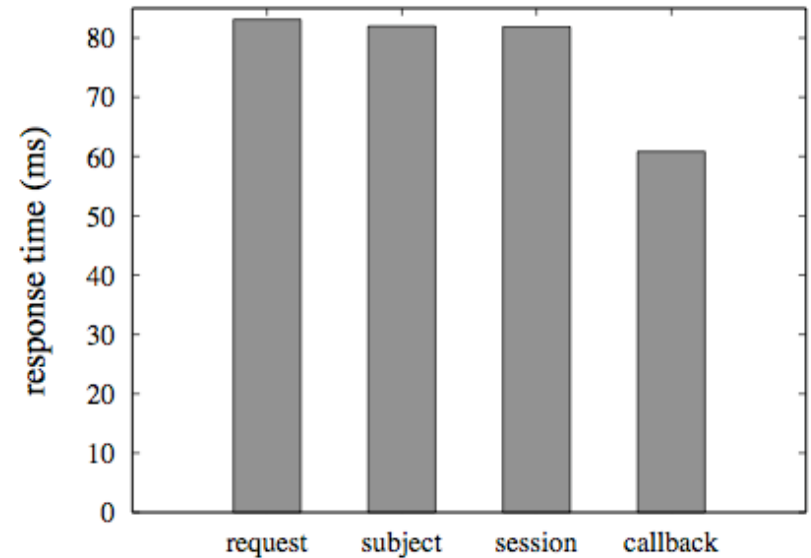
- 500 subjects
- 500 objects
- 3 access rights
- 100 requests/second
- 20 new subjects/second
- 100 active subjects (or sessions)

response time comparison

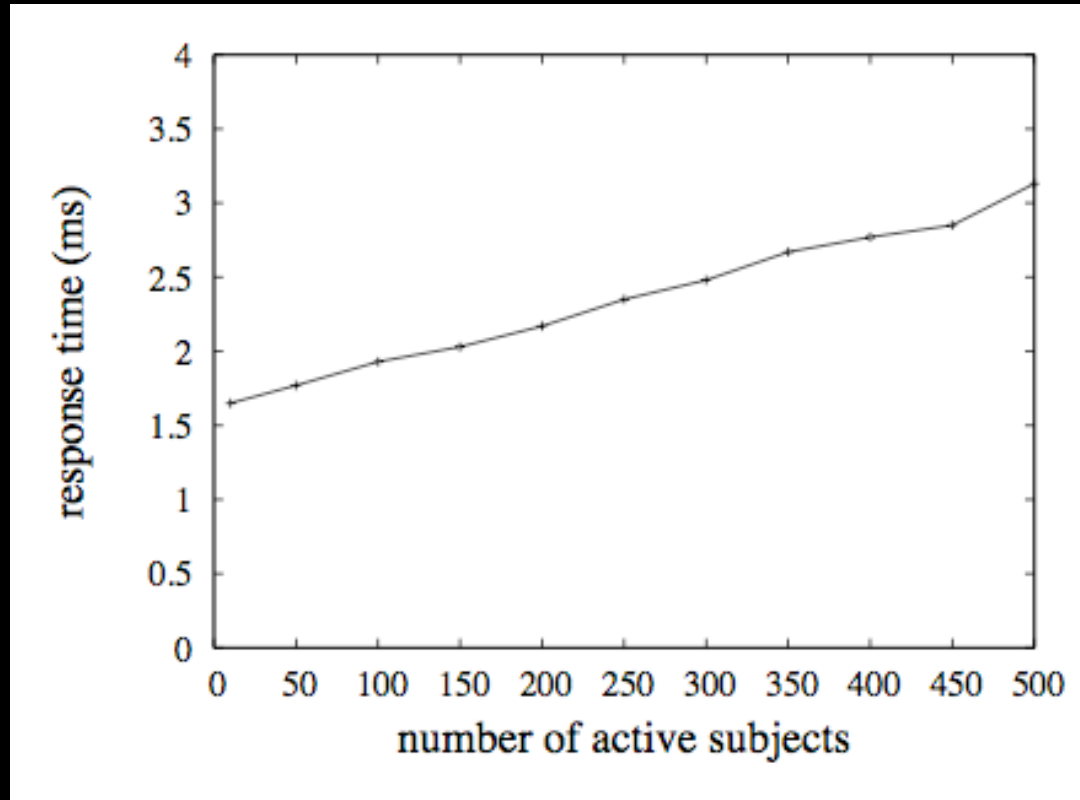
LAN (RTT < 0.1ms)



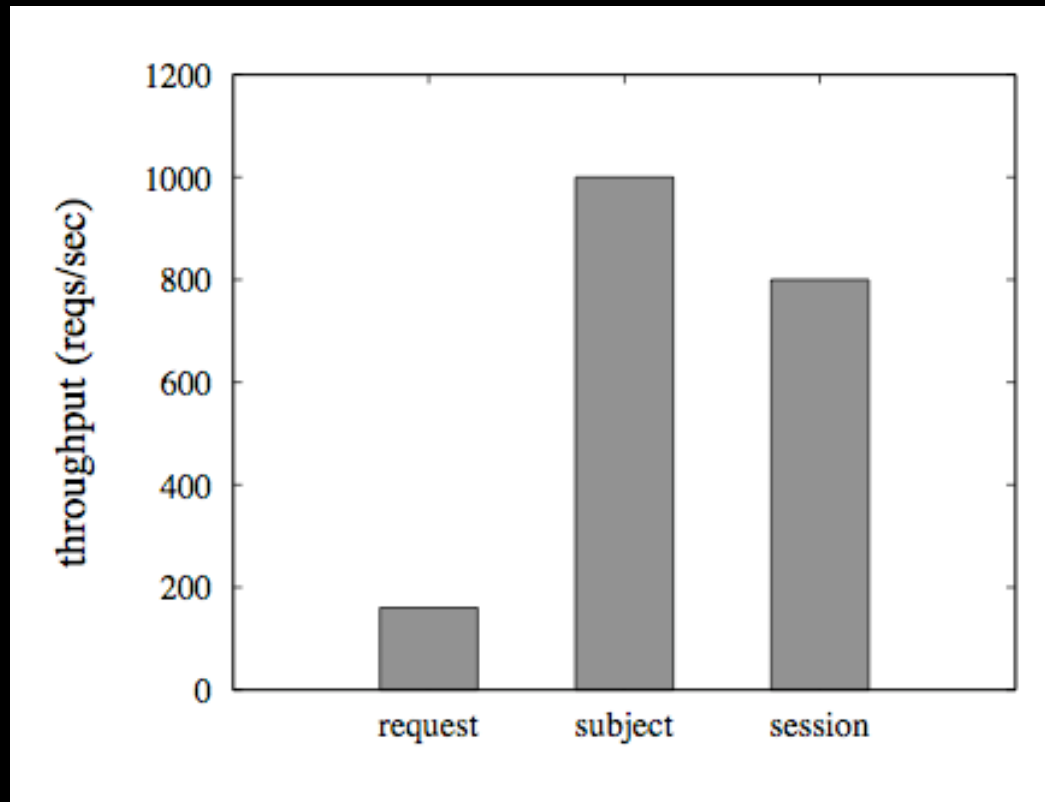
WAN (RTT 40ms)



outstanding subscriptions and latency



subscription frequency and throughput

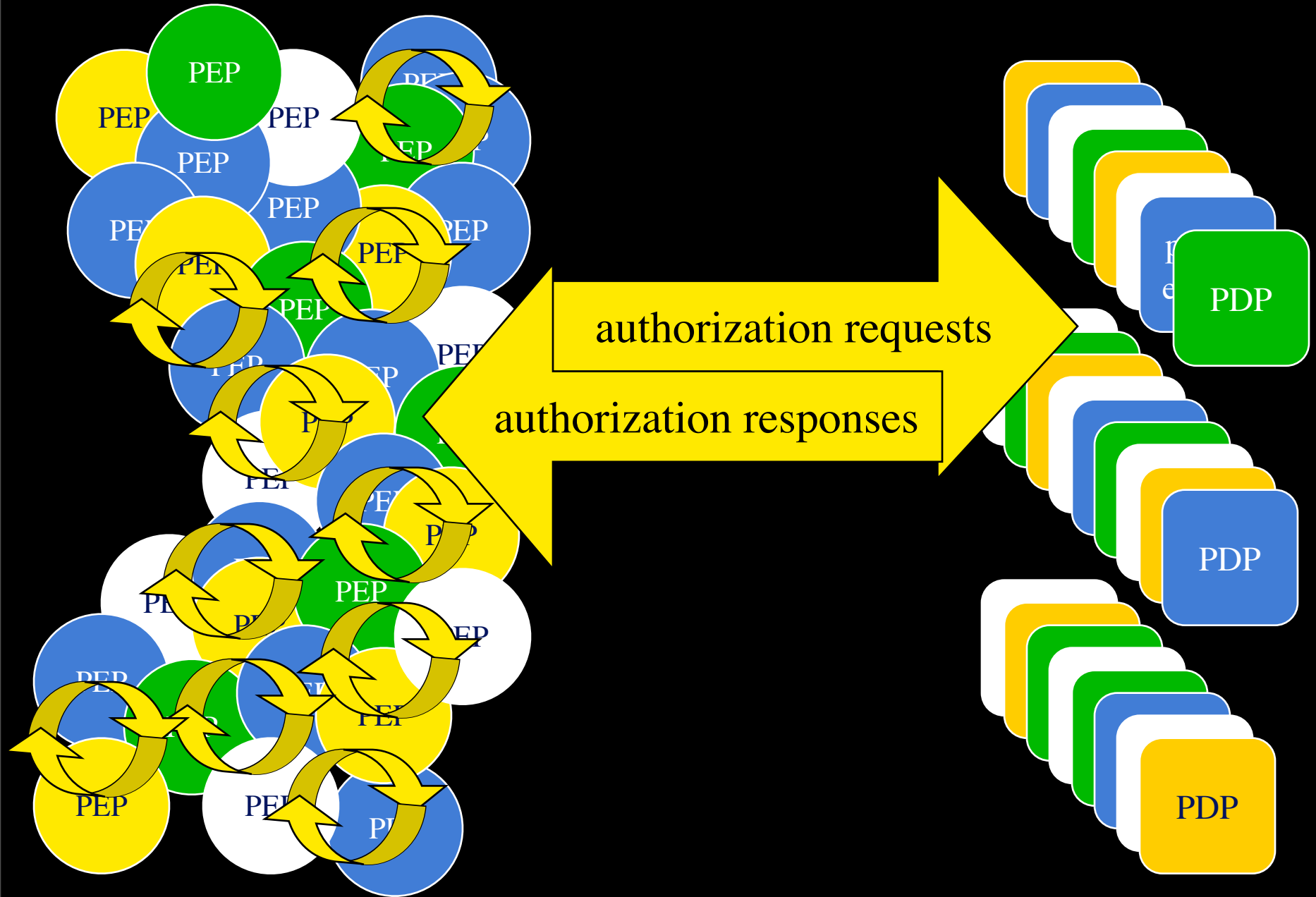


subscription additions/deletions result in ENS matching table updates

conclusions & work in progress

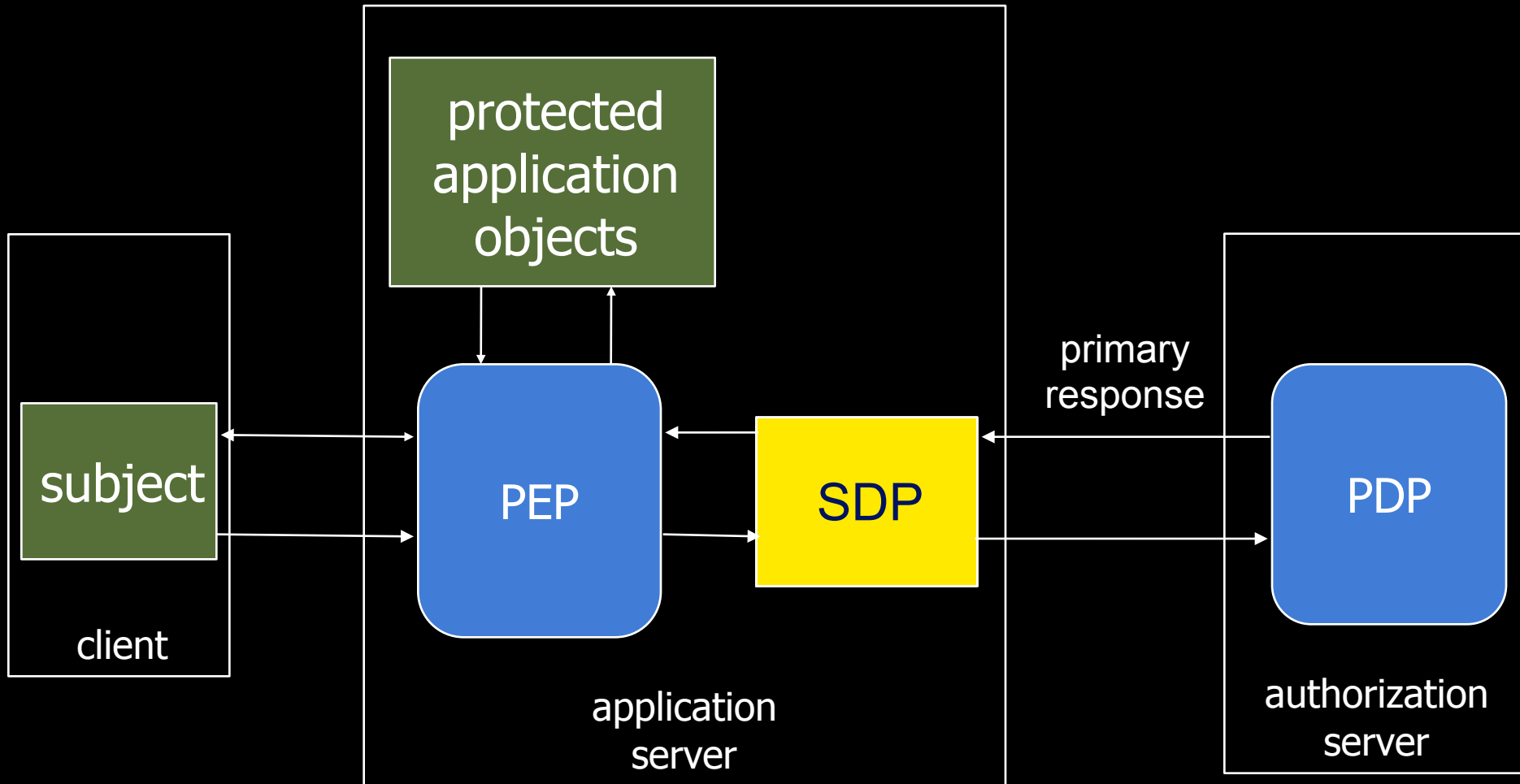
- pub-sub helps to improve system availability
- while employing pub-sub system, aim at
 - low subscription frequency
 - few outstanding subscriptions
- work in progress
 - security of the authorization infrastructure

recycling authorizations

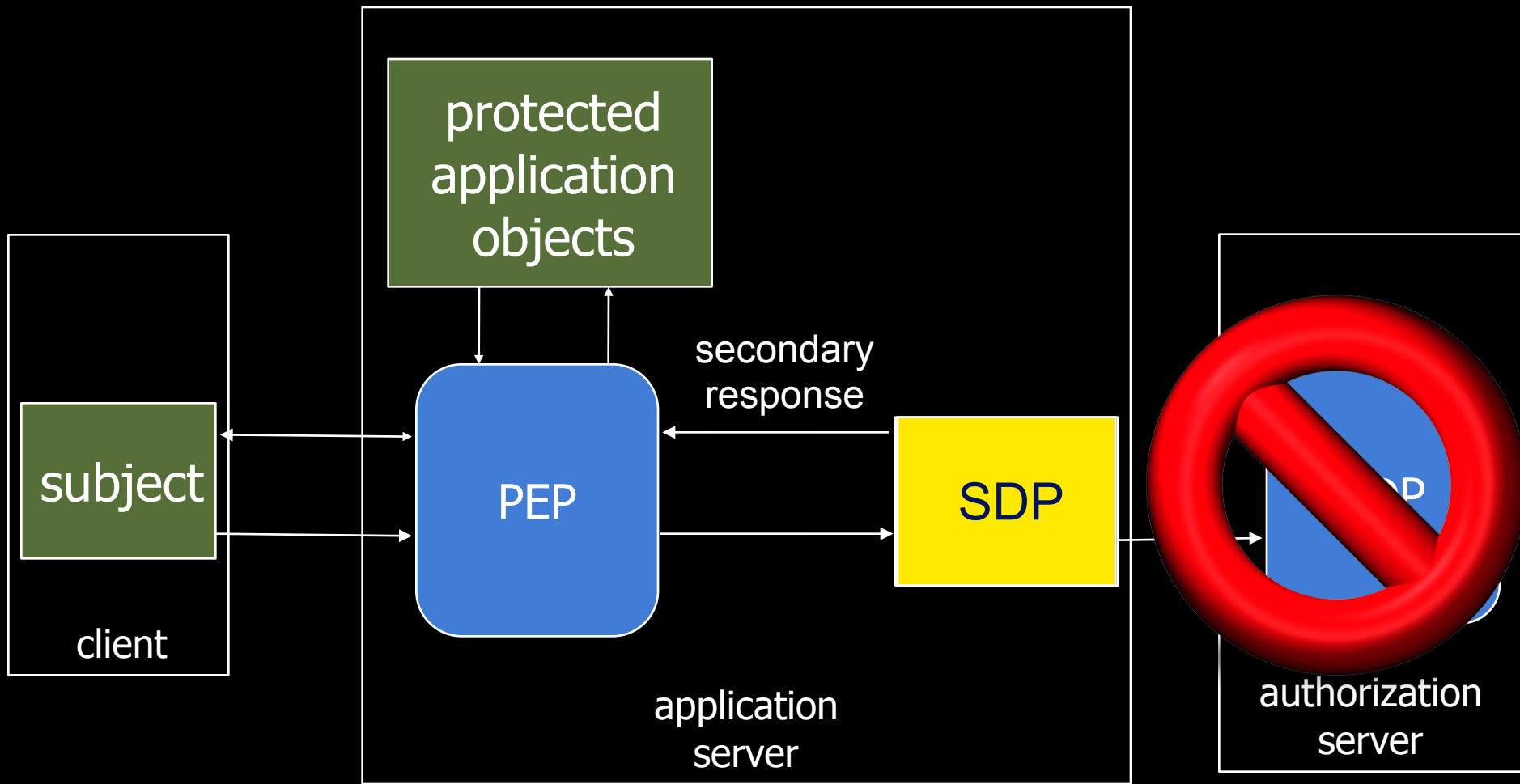


**Secondary and Approximate
Authorization Model
(SAAM)**

what SDP does



what SDP does



SAAM basic elements

▪ request

<subject, object, access right, context, request id>

< s , o , a , c , i >
<{id="Bob", role="customer"}, {id="eB-23"}, view, {date="05-08-15"}, 10 >

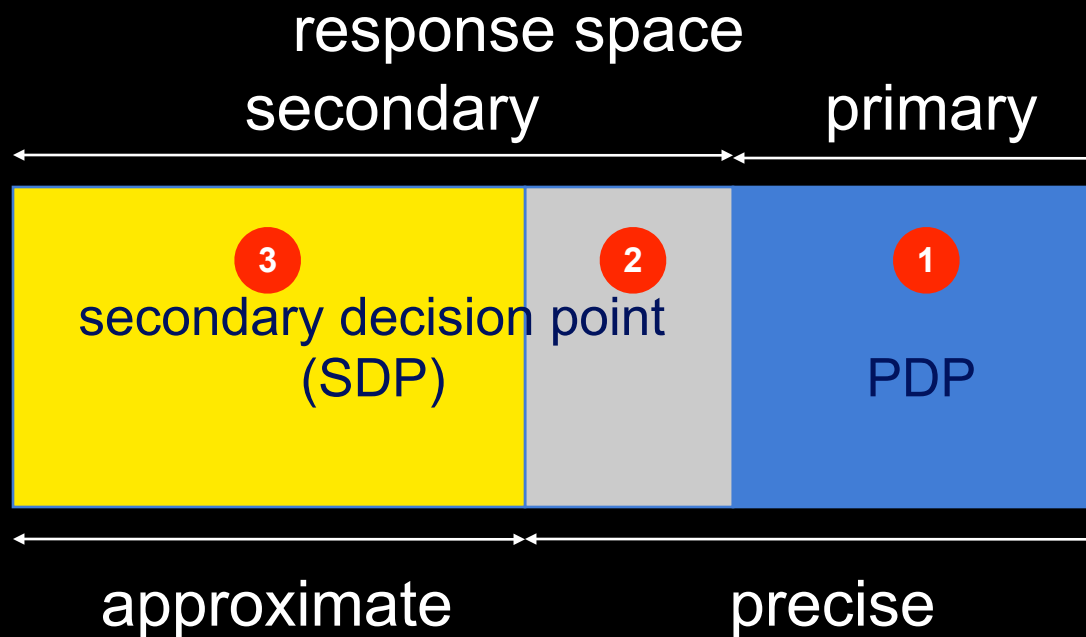
▪ response

<response id, request id, evidence, decision>

< r, i, E, d >
< 1, 10, [], allow >

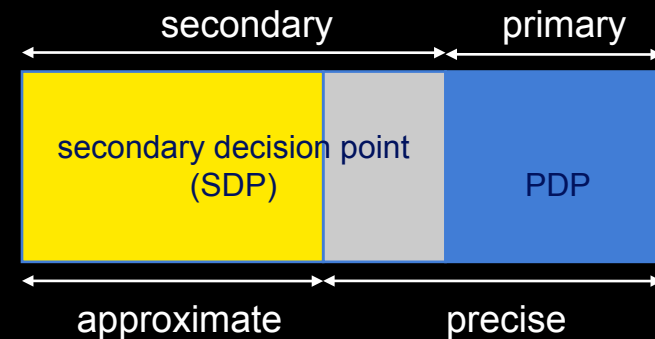
authorization response types

- $\langle \{id="Bob", role="customer"\}, \{id="eB-23"\}, view, \{date="05.06.08"\}, 10 \rangle$ ← equivalent
- ← $\langle 1, 10, [], allow \rangle$ -- primary (from PDP) response
- $\langle \{id="Bob", role="customer"\}, \{id="eB-23"\}, view, \{date="05.06.08"\}, 11 \rangle$
- ← $\langle 2, 11, [1], allow \rangle$ -- secondary and precise response
- $\langle \{id="Alice", role="customer"\}, \{id="eB-23"\}, view, \{date="05.06.08"\}, 12 \rangle$
- ← $\langle 3, 12, [1], allow \rangle$ -- secondary and approximate response

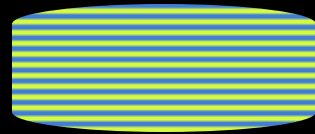


SAAM summary

- basic elements
 - authorization requests $\langle s, o, a, c, i \rangle$
 - authorization responses $\langle r, i, E, d \rangle$
- responses can be
 - primary or secondary
 - precise or approximate
- secondary decision point
 - implemented at PEP
 - uses primary to compute secondary



recycling algorithms



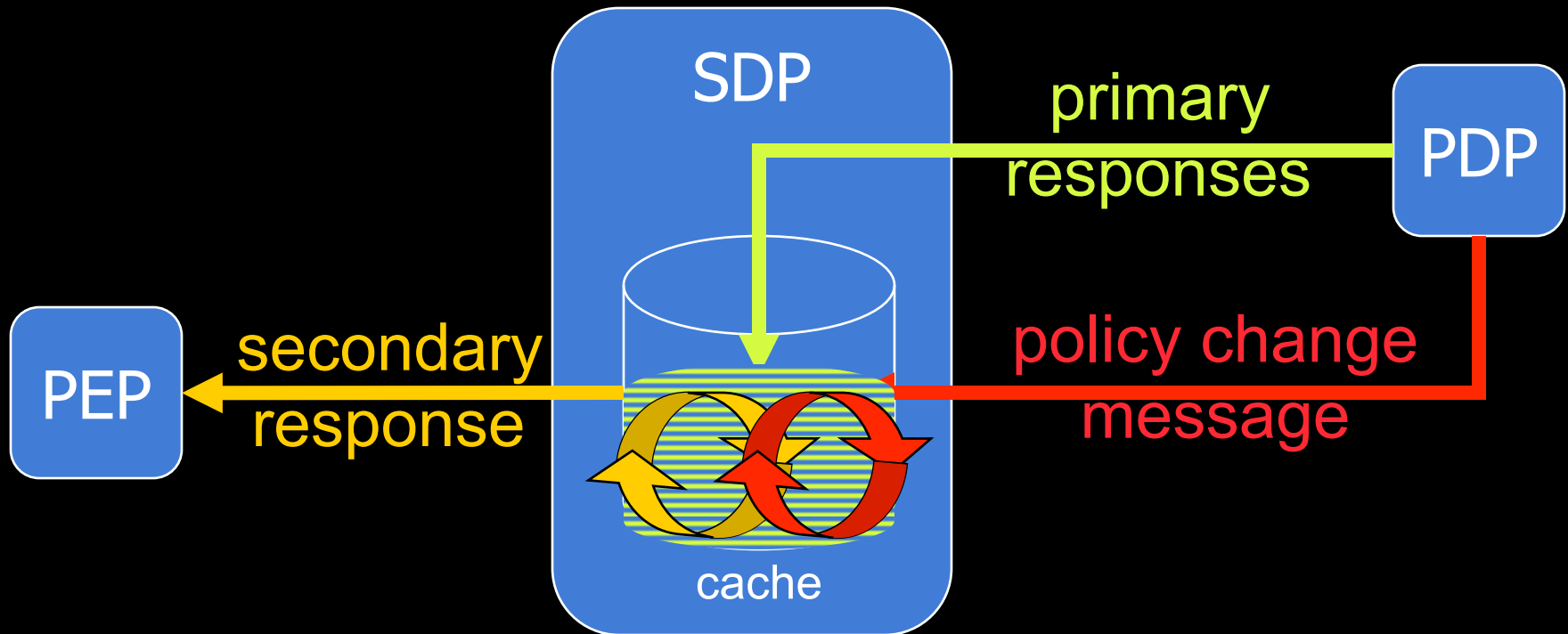
cache construction



decision



cache update

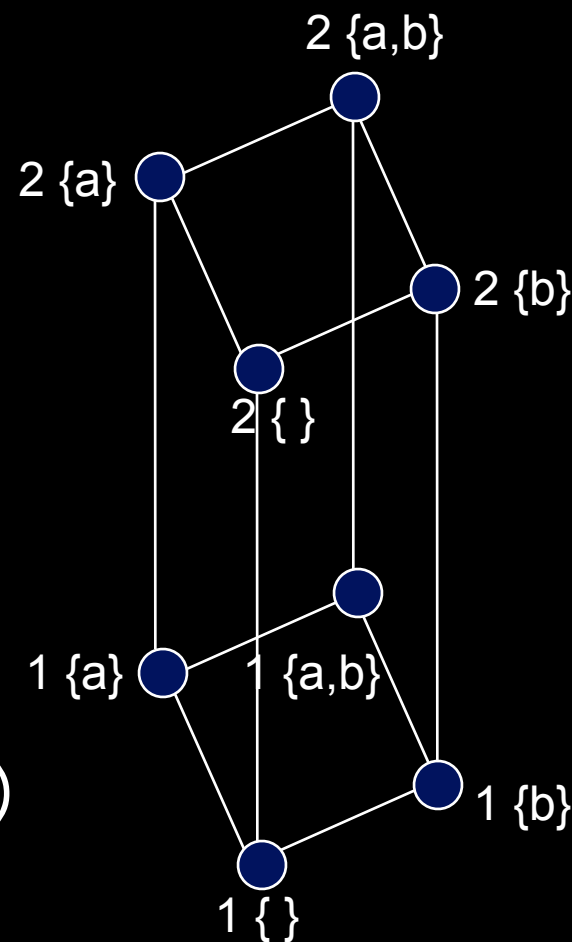


Application of SAAM to Bell LaPadula Policies

J. Crampton, W. Leung, K. Beznosov, “The Secondary and Approximate Authorization Model and its Application to Bell-LaPadula Policies,” in *Proceedings of the ACM Symposium on Access Control Models and Technologies (SACMAT)*, Lake Tahoe, California, USA, 7-9 June, 2006, pp. 111-120.

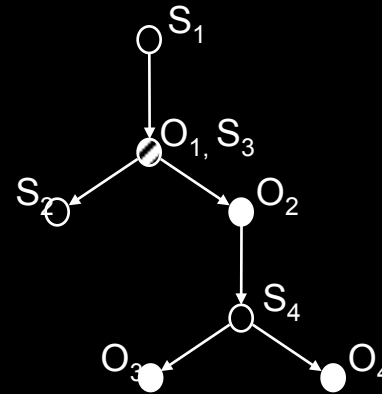
BLP refresher

- S : subjects, O : objects
- DAC
- L : lattice of security labels
- $\lambda: S \cup O \rightarrow L$
- ss-property, *-property:
 - (s, o, read) is allowed $\Rightarrow \lambda(o) \leq \lambda(s)$
 - (s, o, append) is allowed $\Rightarrow \lambda(o) \geq \lambda(s)$
 - (s, o, write) is allowed $\Rightarrow \lambda(o) = \lambda(s)$



What's SAAM_{BLP}?

1. dominance graph (DG)



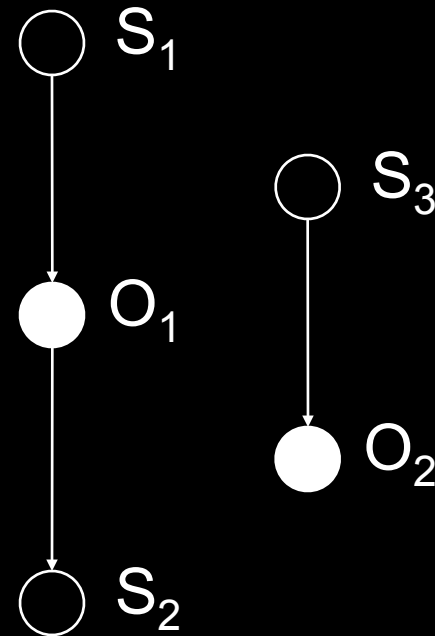
2. algorithms for SDP to

- build DG from primary responses
- compute secondary responses using DG

dominance graph

allow

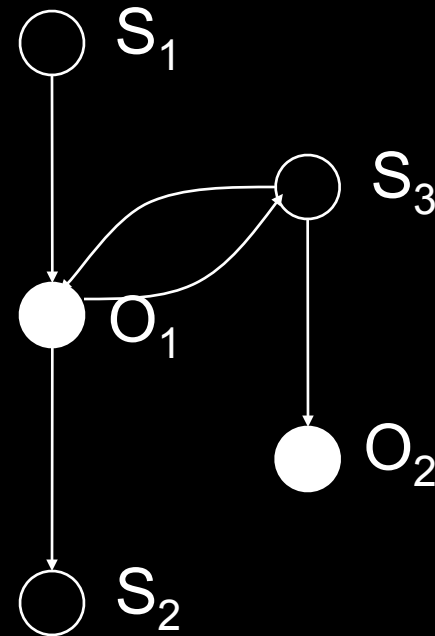
1. (s_1, o_1, read)
2. $(s_2, o_1, \text{append})$
3. (s_3, o_2, read)



dominance graph

allow

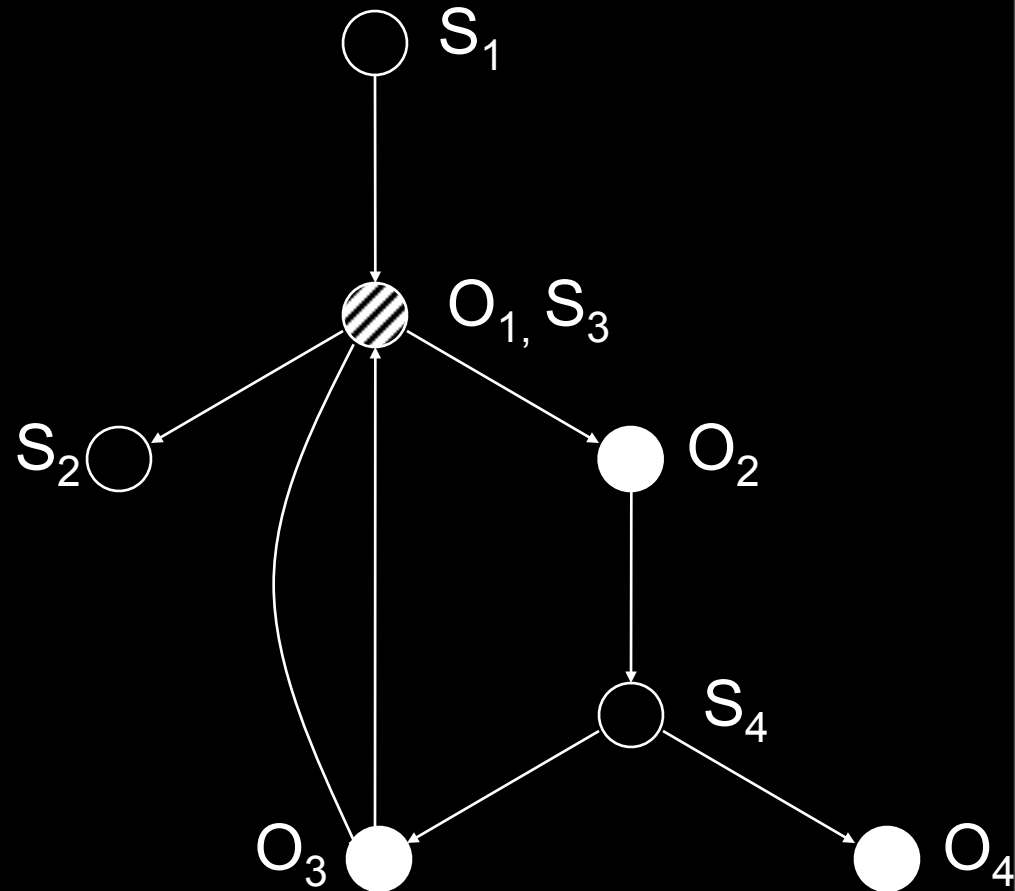
1. (s_1, o_1, read)
2. $(s_2, o_1, \text{append})$
3. (s_3, o_2, read)
4. (s_3, o_1, write)



dominance graph

allow

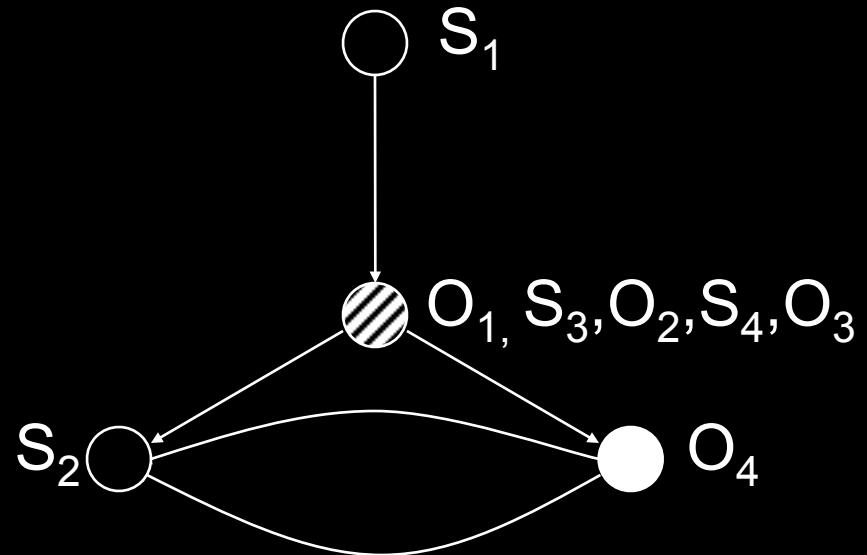
1. (s_1, o_1, read)
2. $(s_2, o_1, \text{append})$
3. (s_3, o_2, read)
4. (s_3, o_1, write)
5. (s_1, o_2, read)
6. $(s_4, o_2, \text{append})$
7. (s_4, o_3, read)
8. (s_4, o_4, read)
9. (s_3, o_3, write)



dominance graph

allow

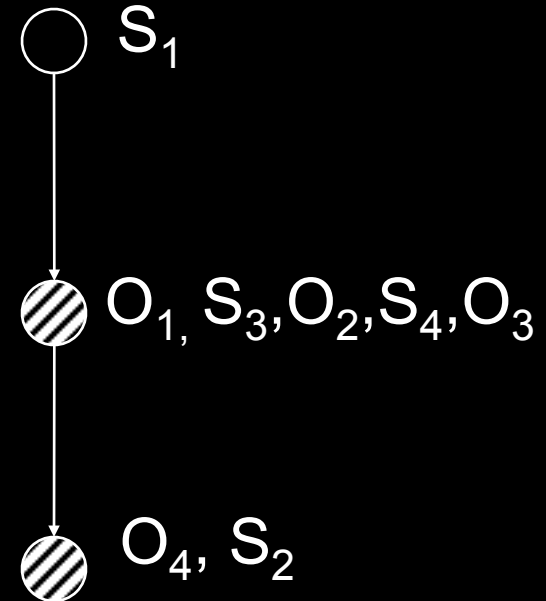
1. (s_1, o_1, read)
2. $(s_2, o_1, \text{append})$
3. (s_3, o_2, read)
4. (s_3, o_1, write)
5. (s_1, o_2, read)
6. $(s_4, o_2, \text{append})$
7. (s_4, o_3, read)
8. (s_4, o_4, read)
9. (s_3, o_3, write)
10. (s_2, o_4, write)



dominance graph

allow

1. (s_1, o_1, read)
2. $(s_2, o_1, \text{append})$
3. (s_3, o_2, read)
4. (s_3, o_1, write)
5. (s_1, o_2, read)
6. $(s_4, o_2, \text{append})$
7. (s_4, o_3, read)
8. (s_4, o_4, read)
9. (s_3, o_3, write)
10. (s_2, o_4, write)



SDP may allow:

- (S_1, O_4, read)
- (S_4, O_1, write)
- $(S_2, O_3, \text{append})$

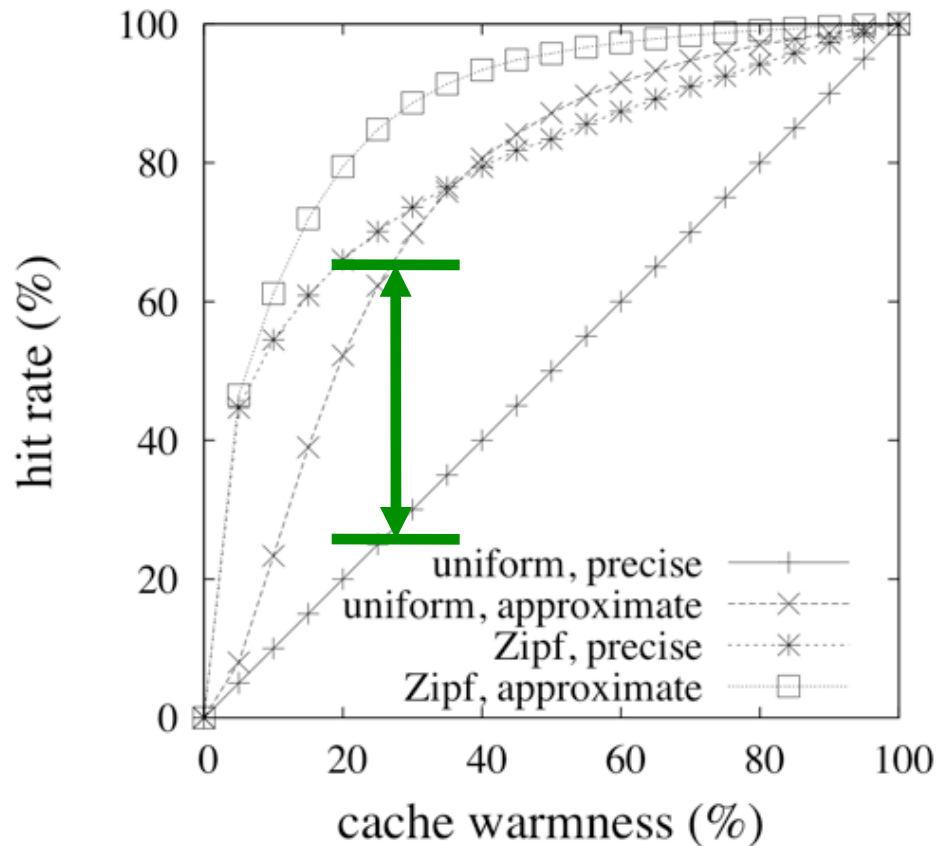
SDP cannot decide:

- (S_2, O_3, read)
- (S_1, O_4, write)
- $(S_1, O_4, \text{append})$

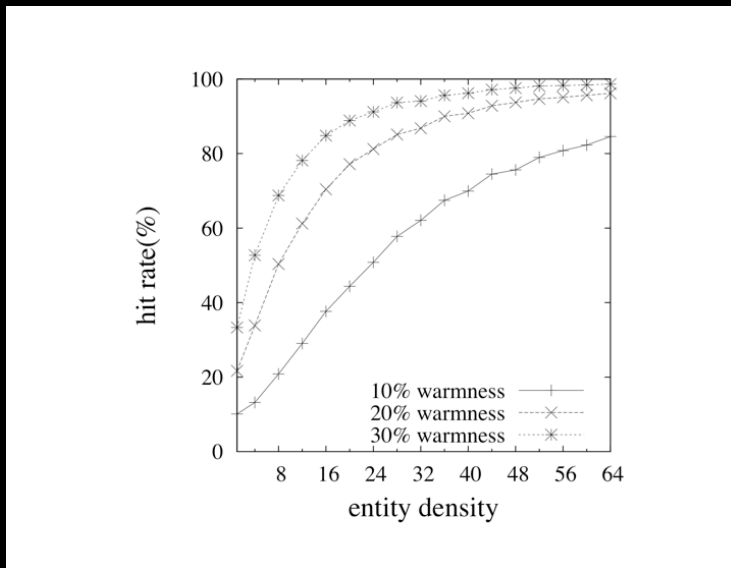
SAAM_{BLP} evaluation

hit rate

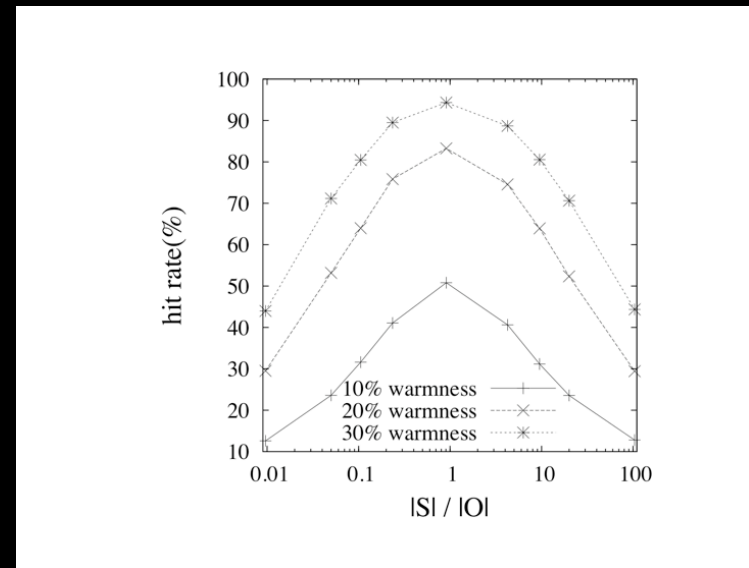
BLP policy: 5 levels, 5 categories, 50 subjects, 1,000 objects, 2 rights



impact of various system parameters



density of subjects and objects in the lattice



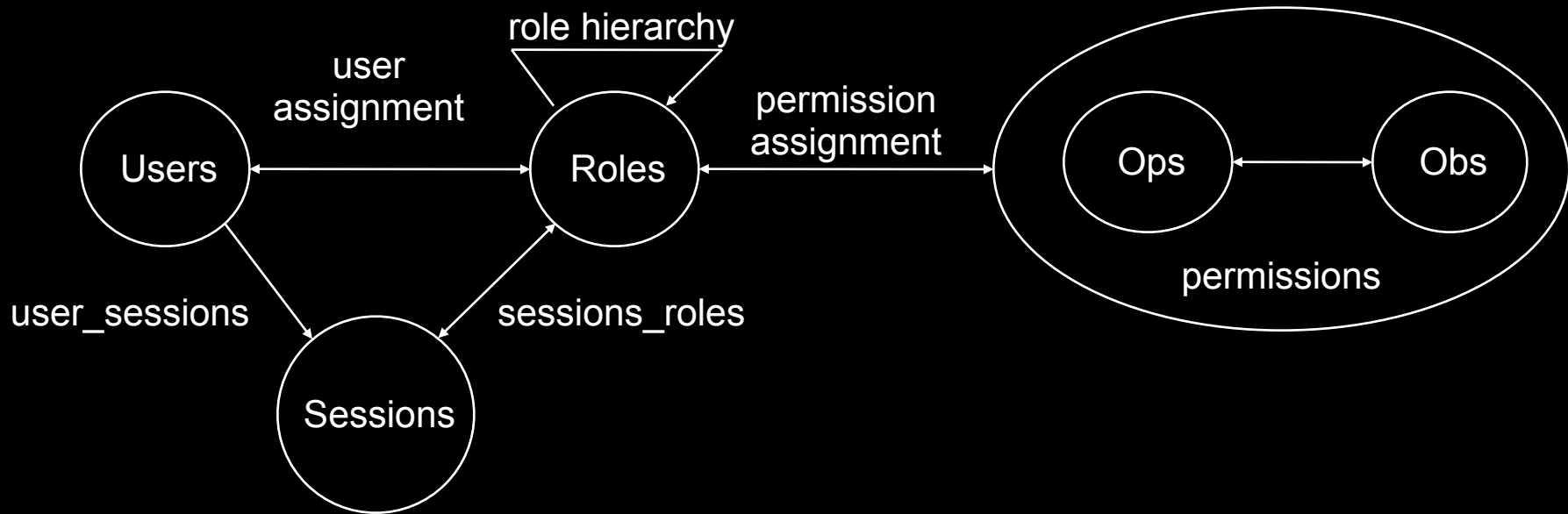
subject/object ratio

J. Crampton, W. Leung, K. Beznosov, "The Secondary and Approximate Authorization Model and its Application to Bell-LaPadula Policies," in *Proceedings of the ACM Symposium on Access Control Models and Technologies (SACMAT)*, Lake Tahoe, California, USA, 7-9 June, 2006, pp. 111-120.

SAAM_{RBAC}: SAAM for RBAC

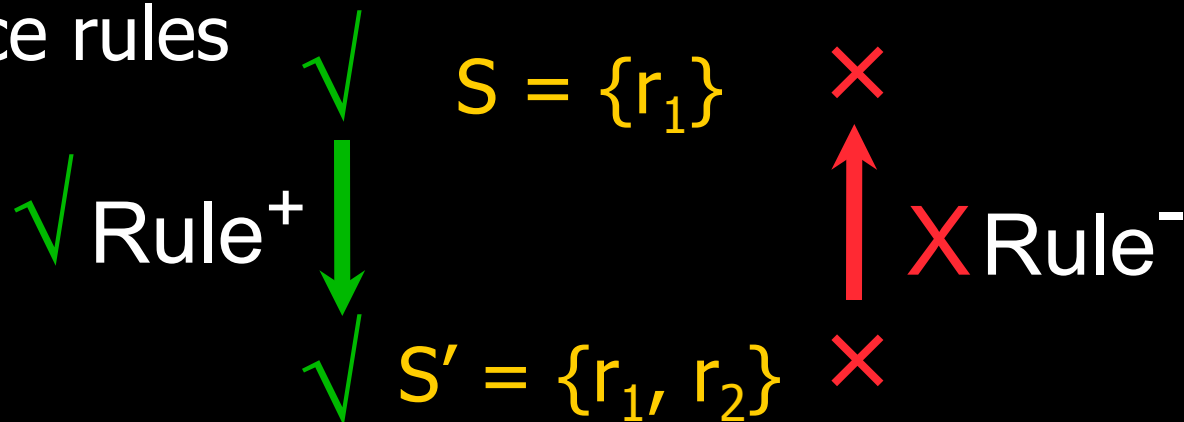
Q. Wei, J. Crampton, K. Beznosov, M. Ripeanu, “[Authorization Recycling in RBAC Systems](#)” to appear in Proceedings of the ACM Symposium on Access Control Models and Technologies (SACMAT), Estes Park, Colorado, 11-13 June 2008.

RBAC review



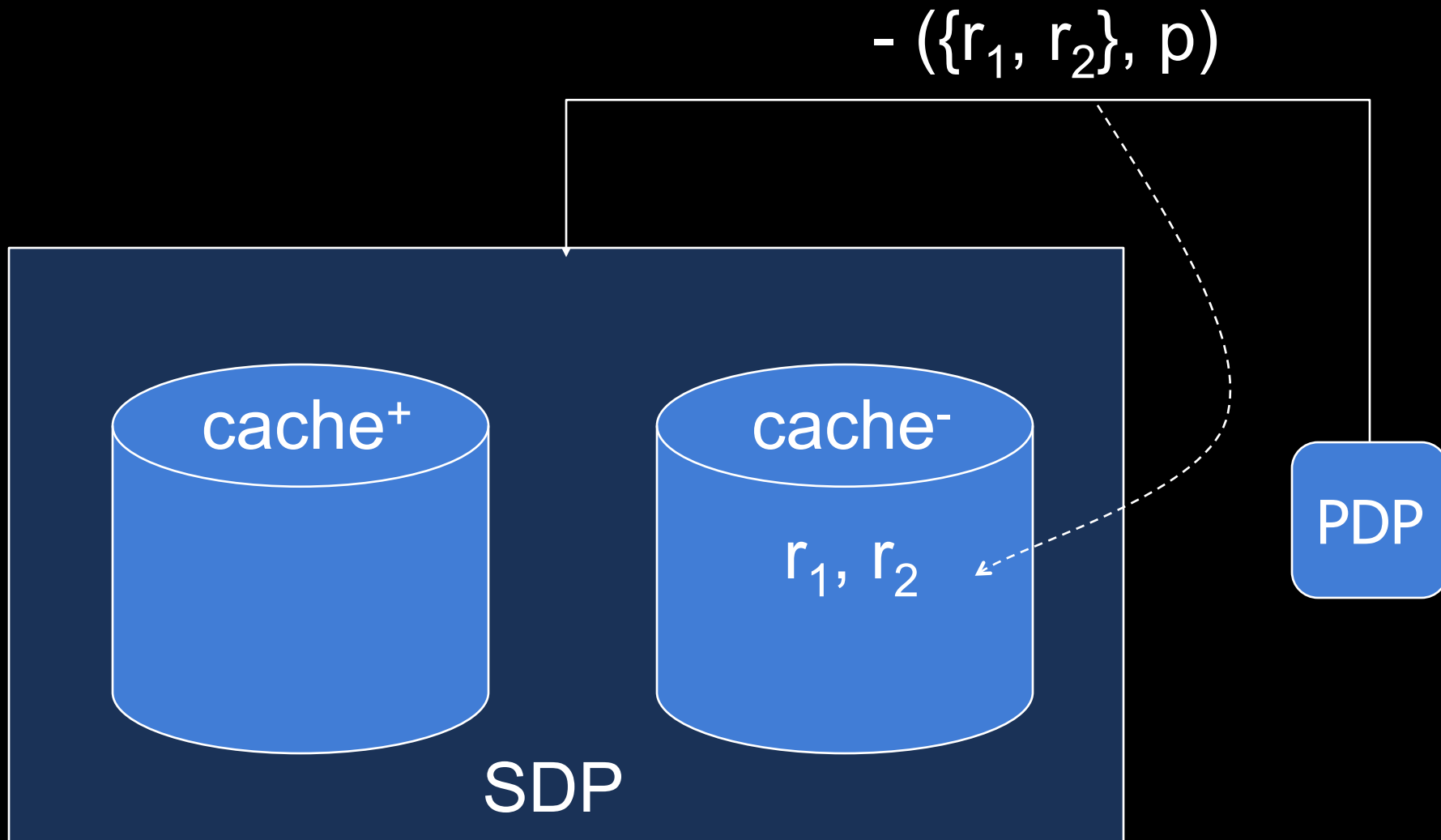
preliminaries

- request: issued by a subject for a permission.
 - $\text{request}=(s,p)$
- \pm : denotes the decision to a request.
 - $\text{response}=+(s,p)$ or $-(s,p)$
- subject: modeled as a set of roles.
 - $s= \{r_2, r_3, r_4\}$
- inference rules

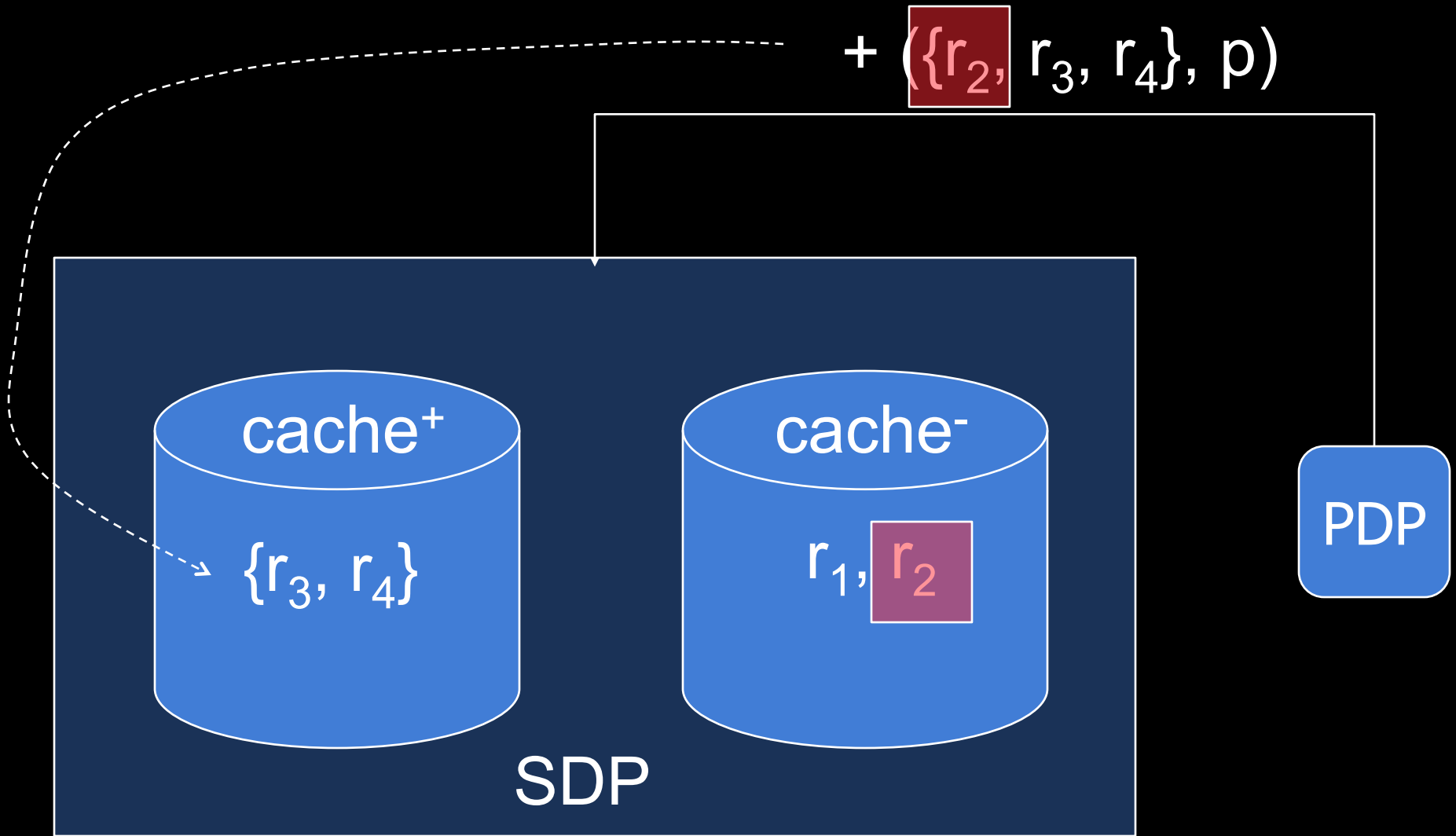


example

caching first negative decision

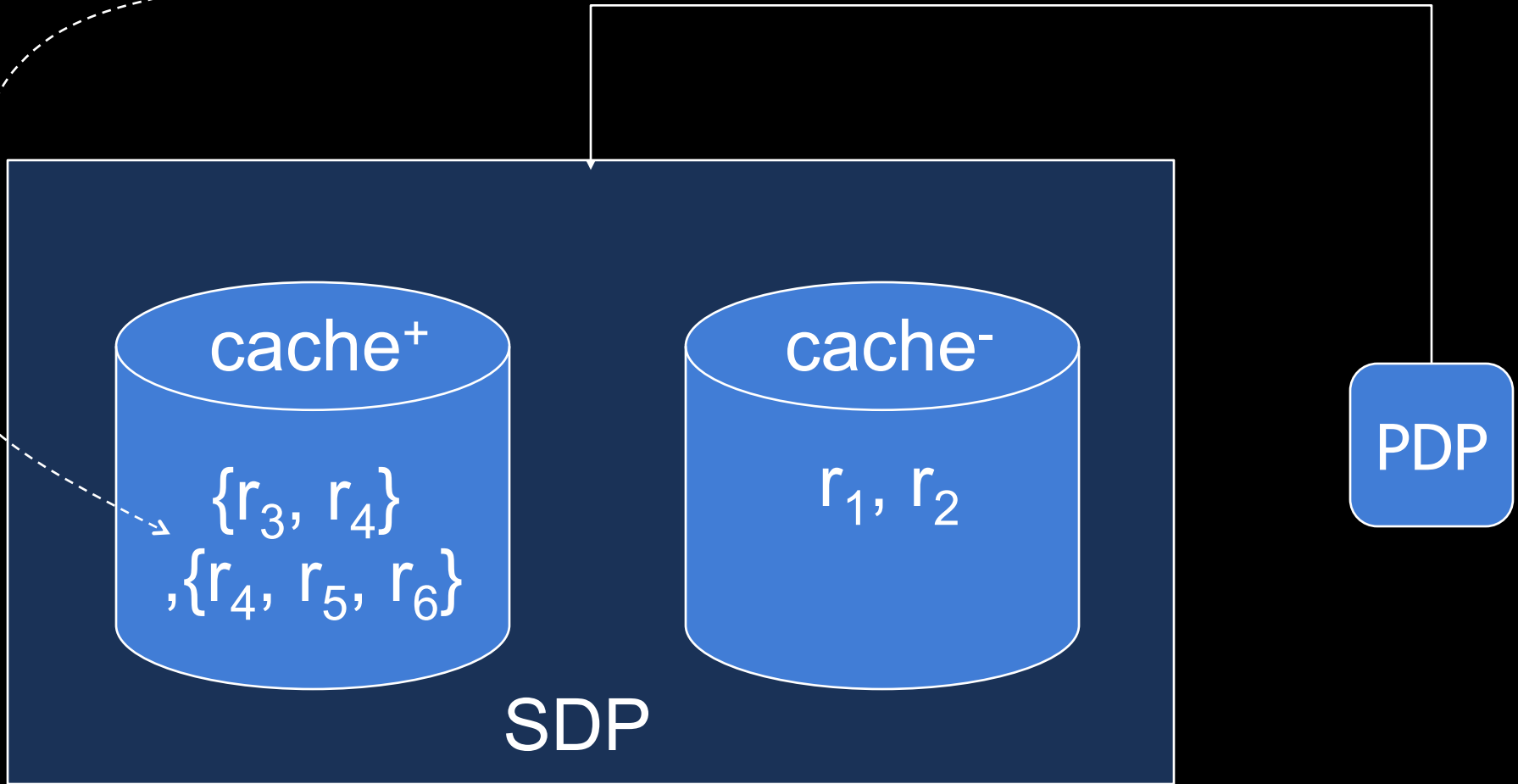


caching first positive decision



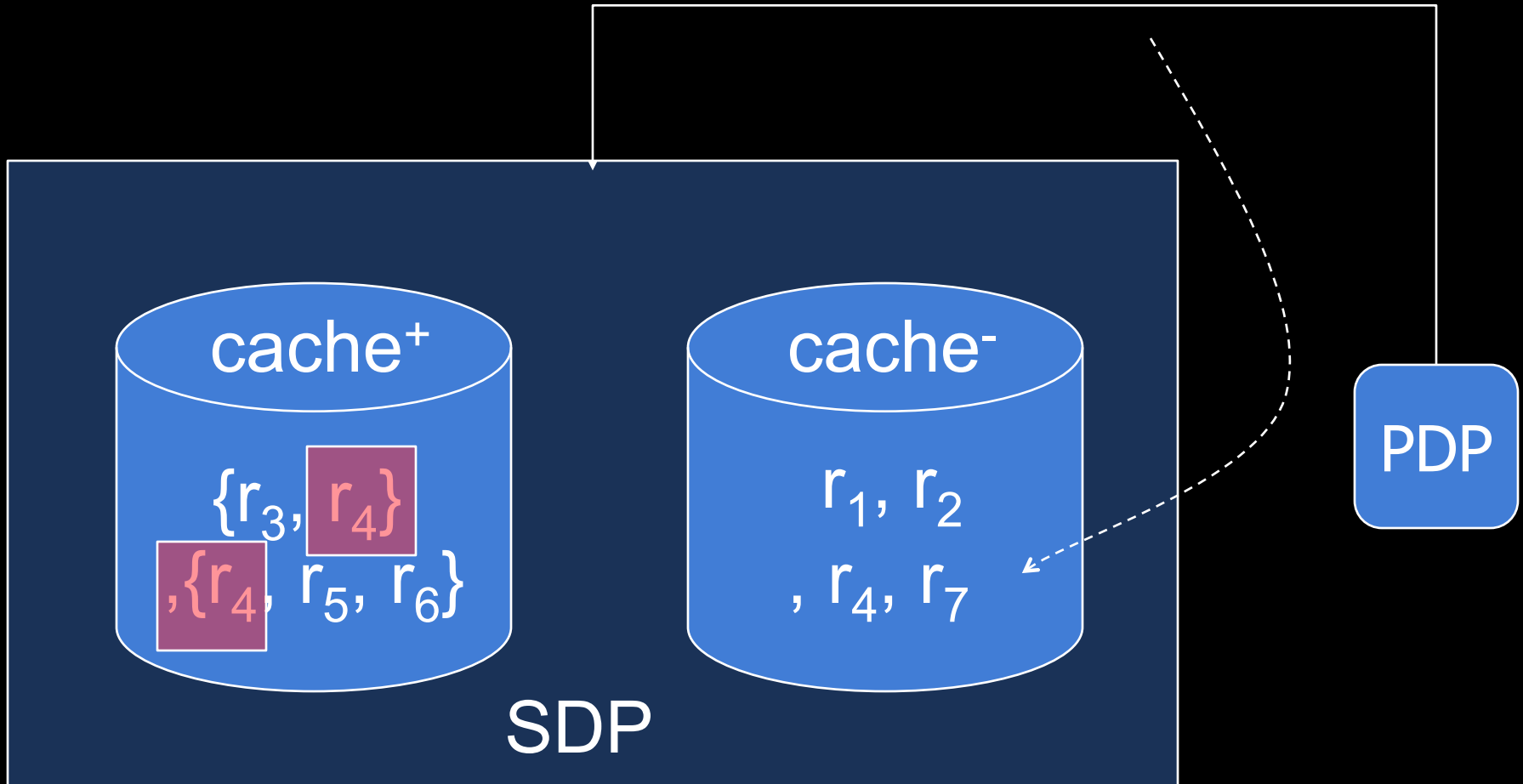
caching second positive decision

+ $(\{r_4, r_5, r_6\}, p)$



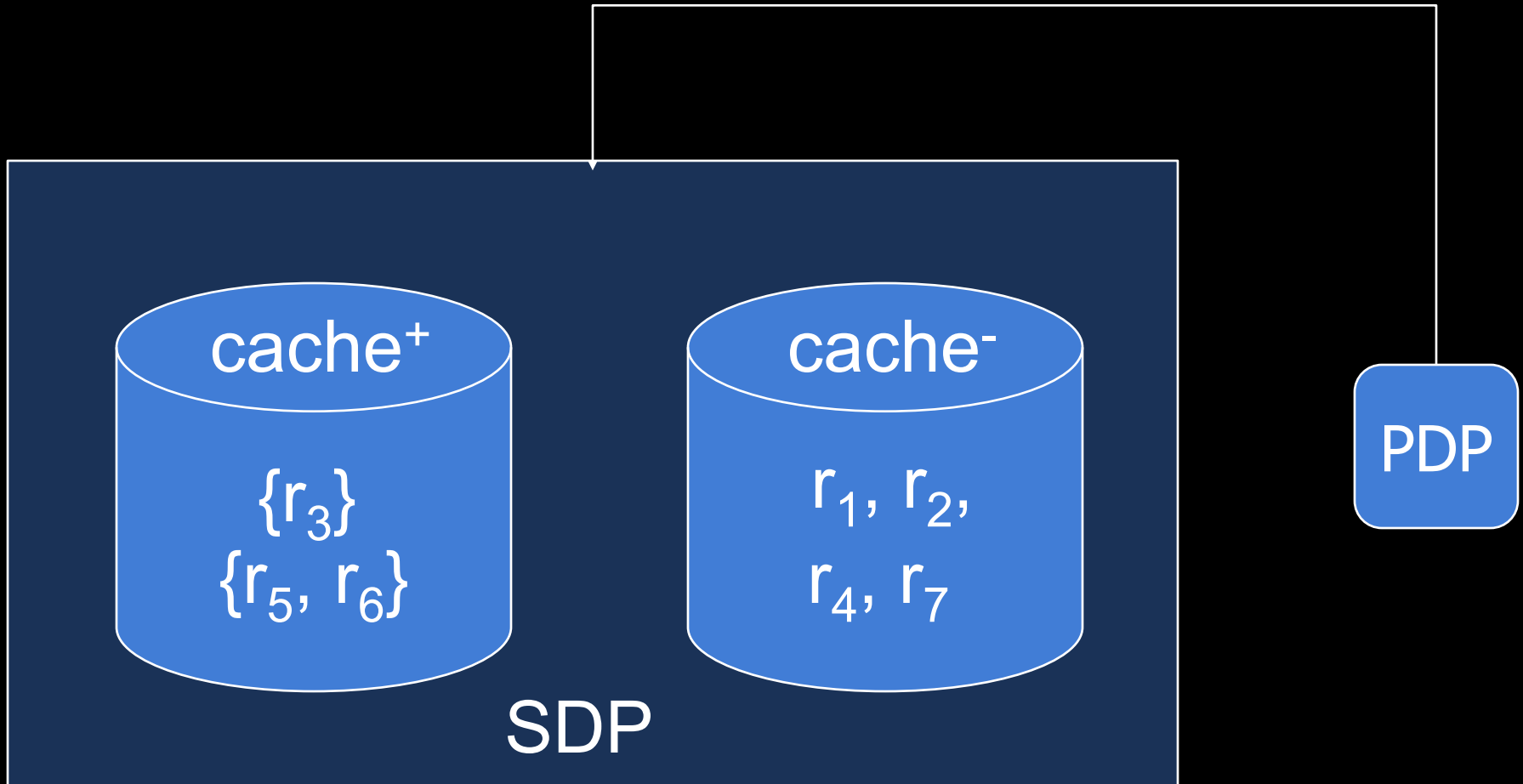
caching second negative decision

- $(\{r_4, r_7\}, p)$

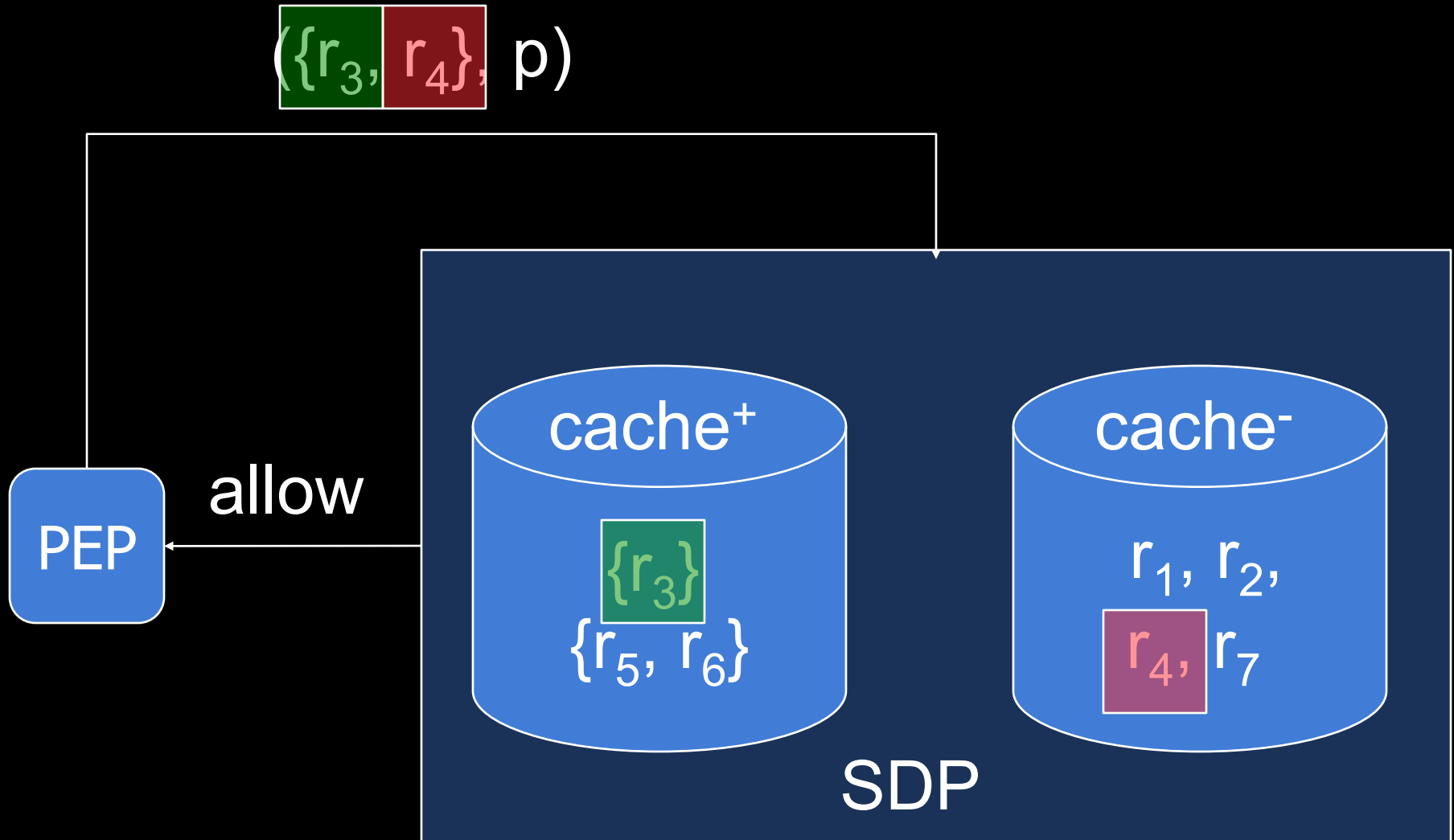


caching second negative decision

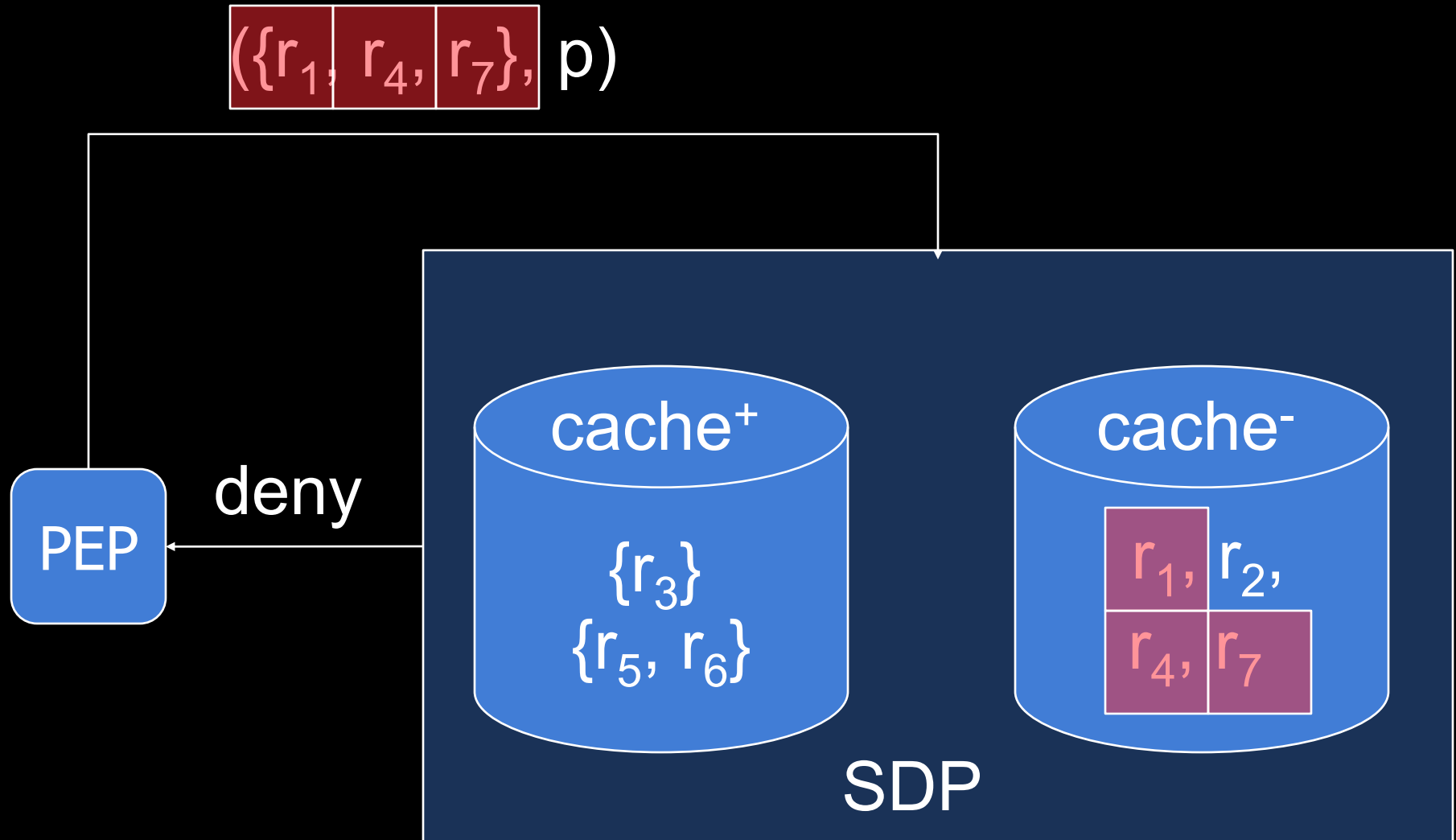
- ($\{r_4, r_7\}, p$)



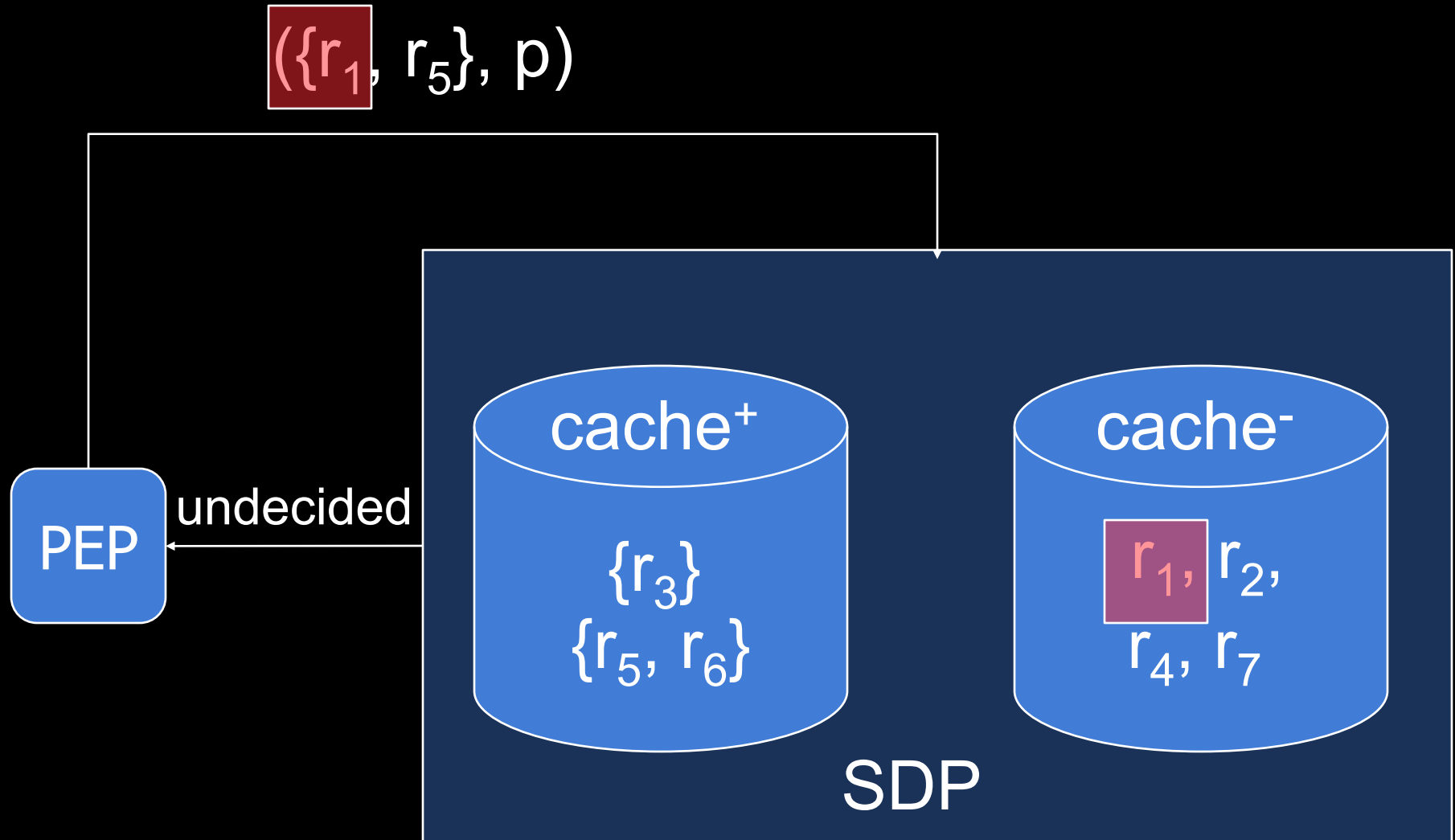
computing allowing authorization



computing denying authorization



computing undecided authorization

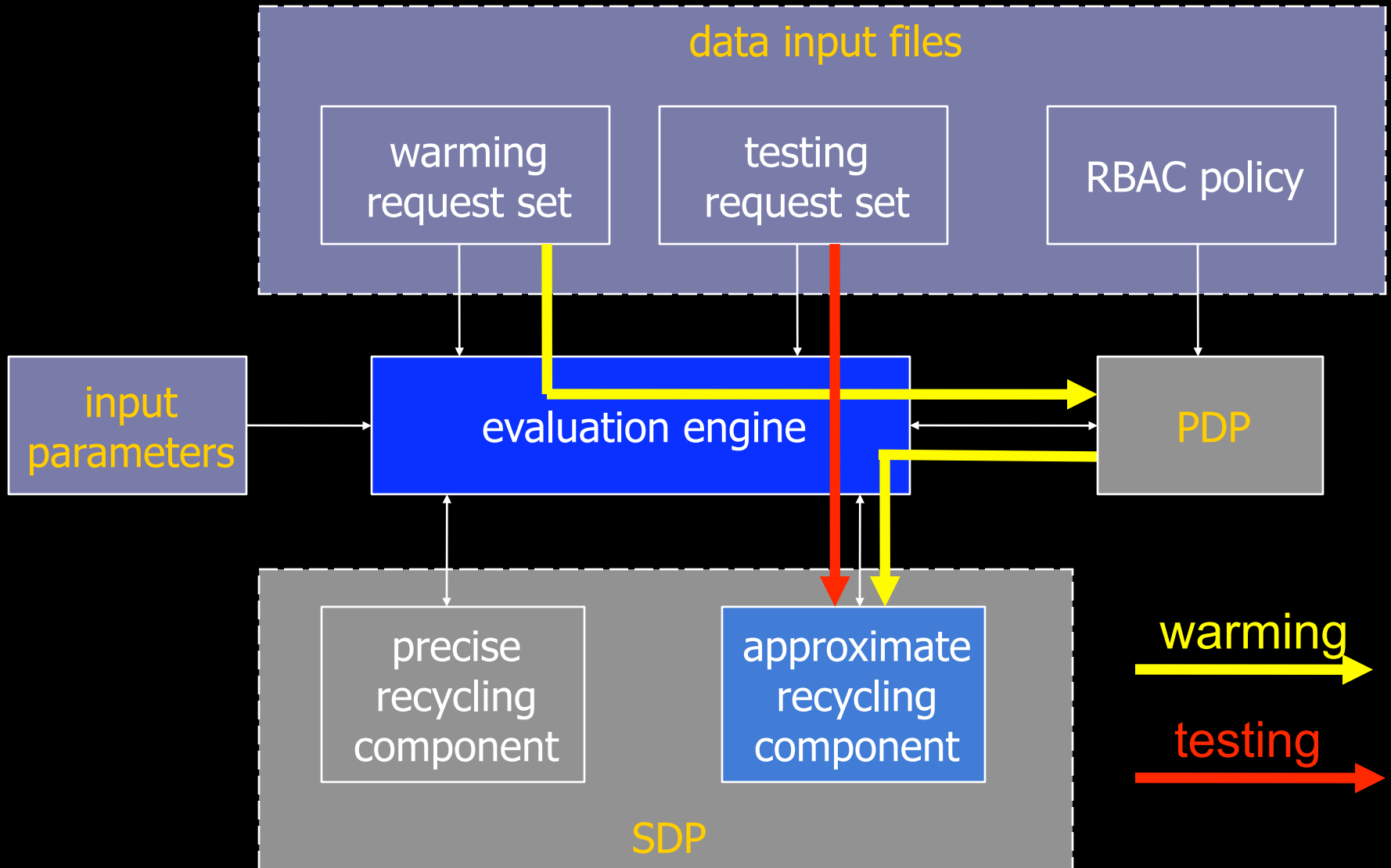


SAAM_{RBAC} evaluation

evaluation metrics

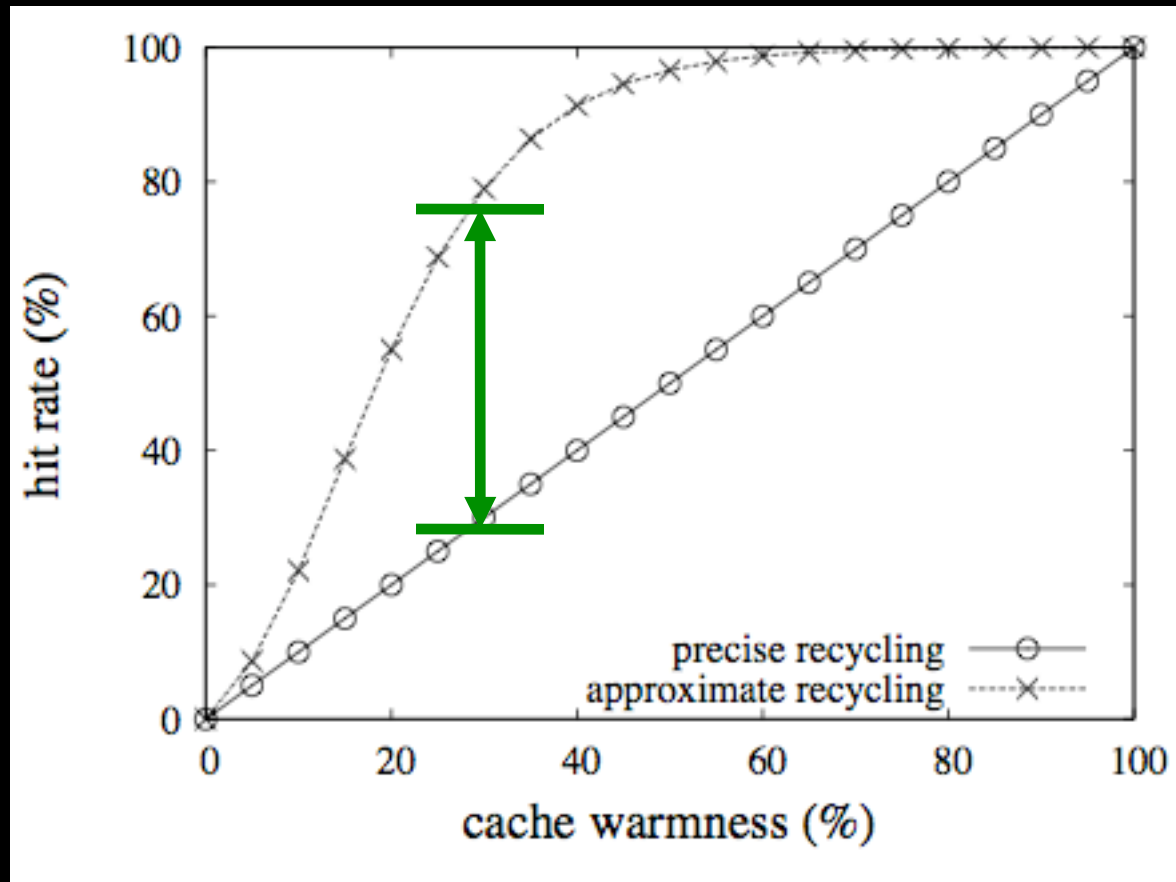
- SDP hit rate
- SDP inference time
 - the time used to infer approximate responses
 - less inference time, more efficient the system

evaluation methodology

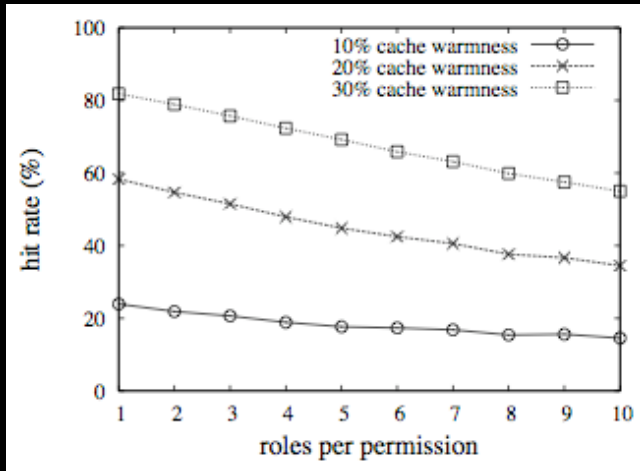


hit rate

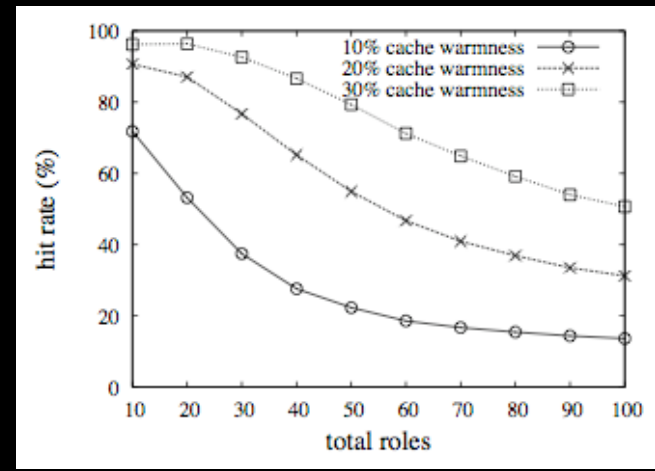
RBAC policy: 100 subjects, 1,000 objects, 50 roles
uniform distribution



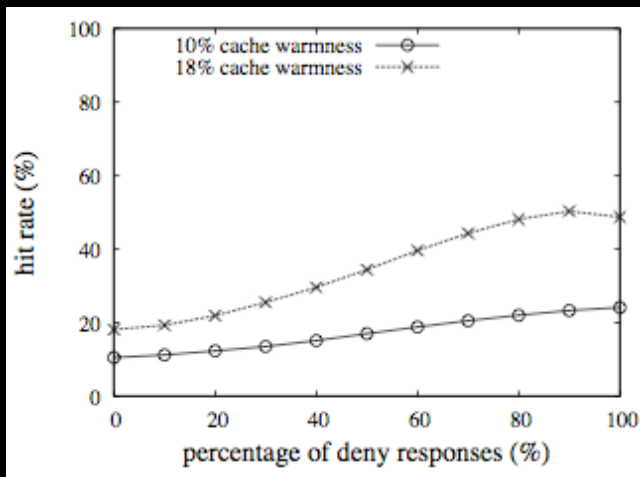
impact of various system parameters



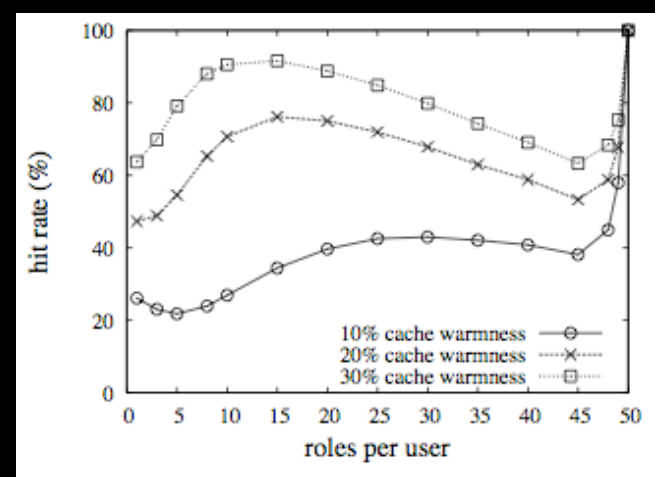
roles per permission



total roles

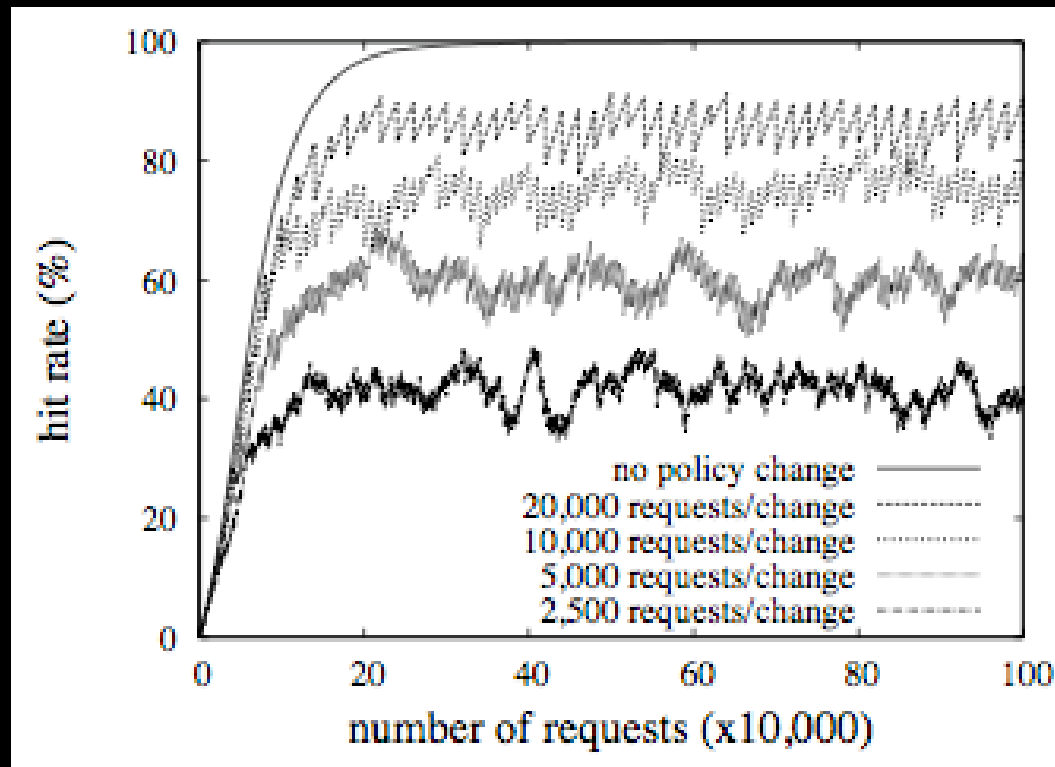


deny vs. allow responses

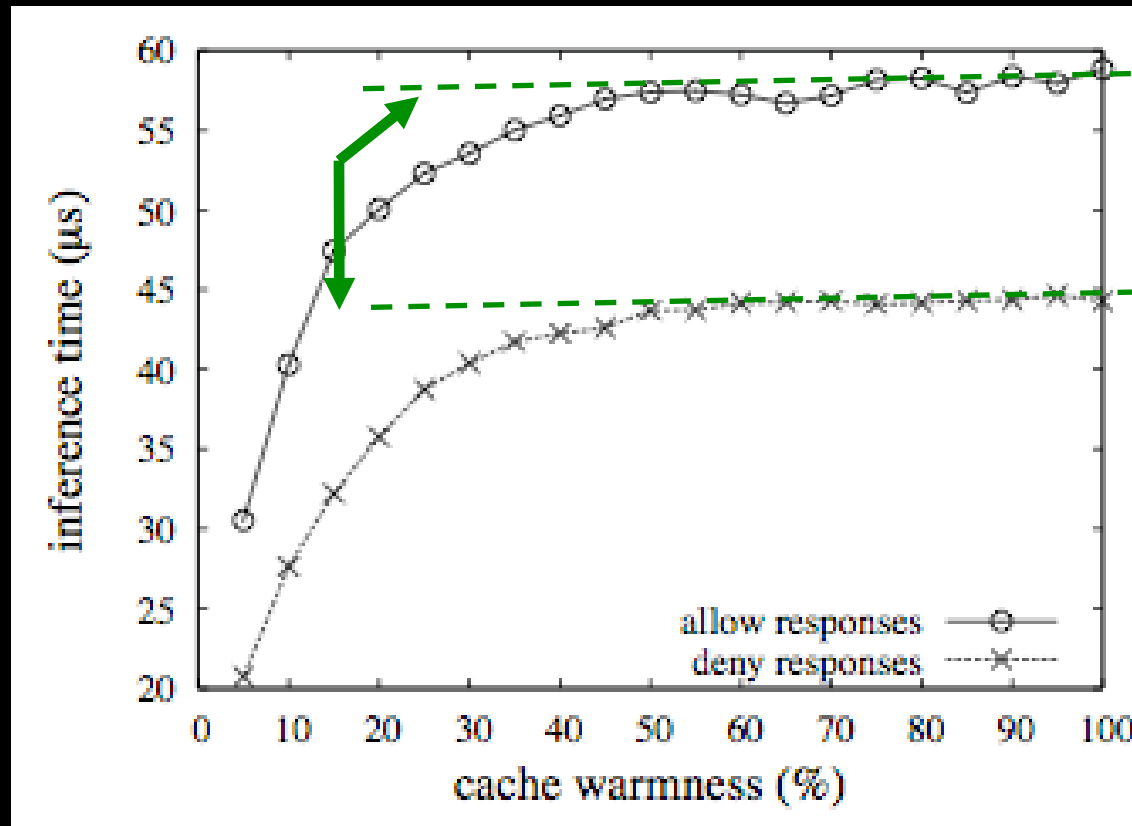


roles per user

impact of policy changes



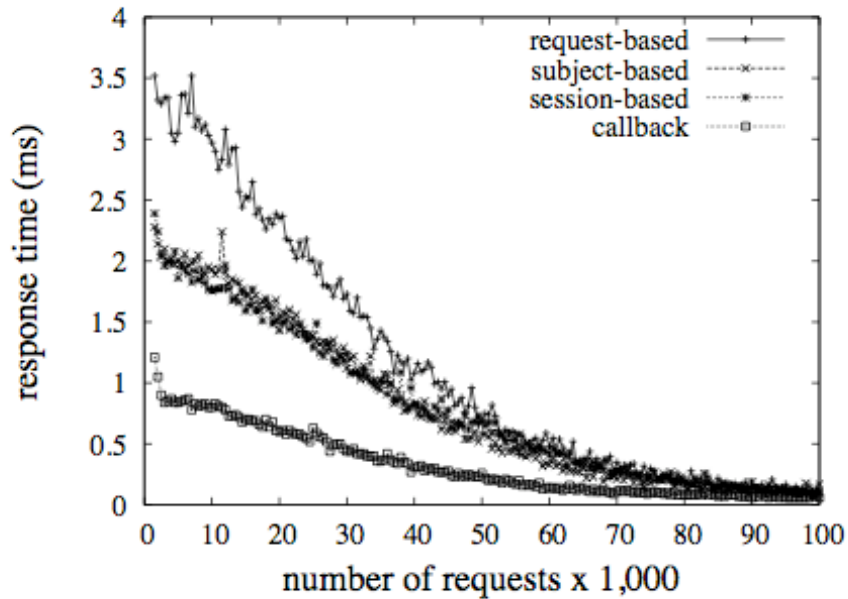
inference time



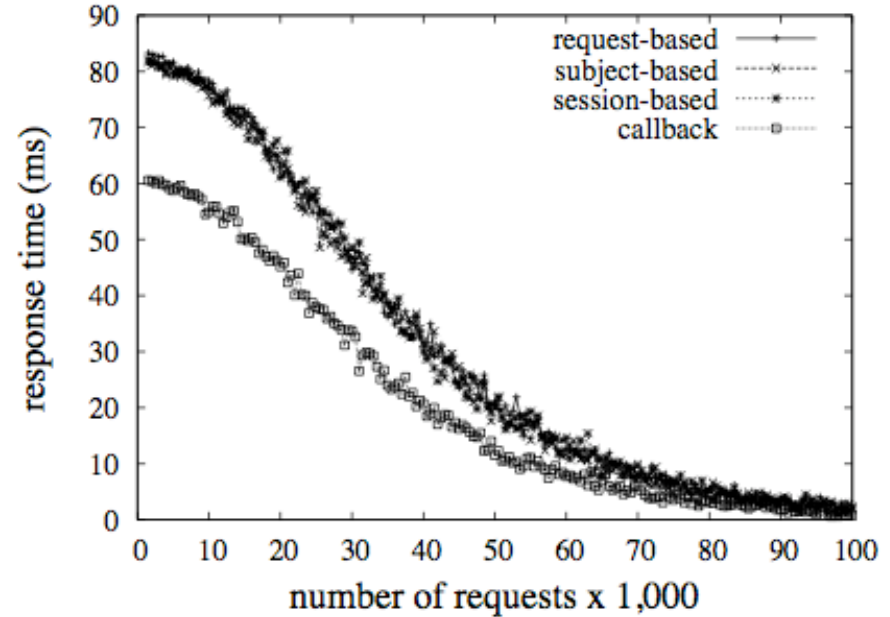
Q. Wei, J. Crampton, K. Beznosov, M. Ripeanu, “**Authorization Recycling in RBAC Systems**” to appear in Proceedings of the ACM Symposium on Access Control Models and Technologies (SACMAT), Estes Park, Colorado, 11-13 June 2008.

combining pub-sub & recycling

LAN

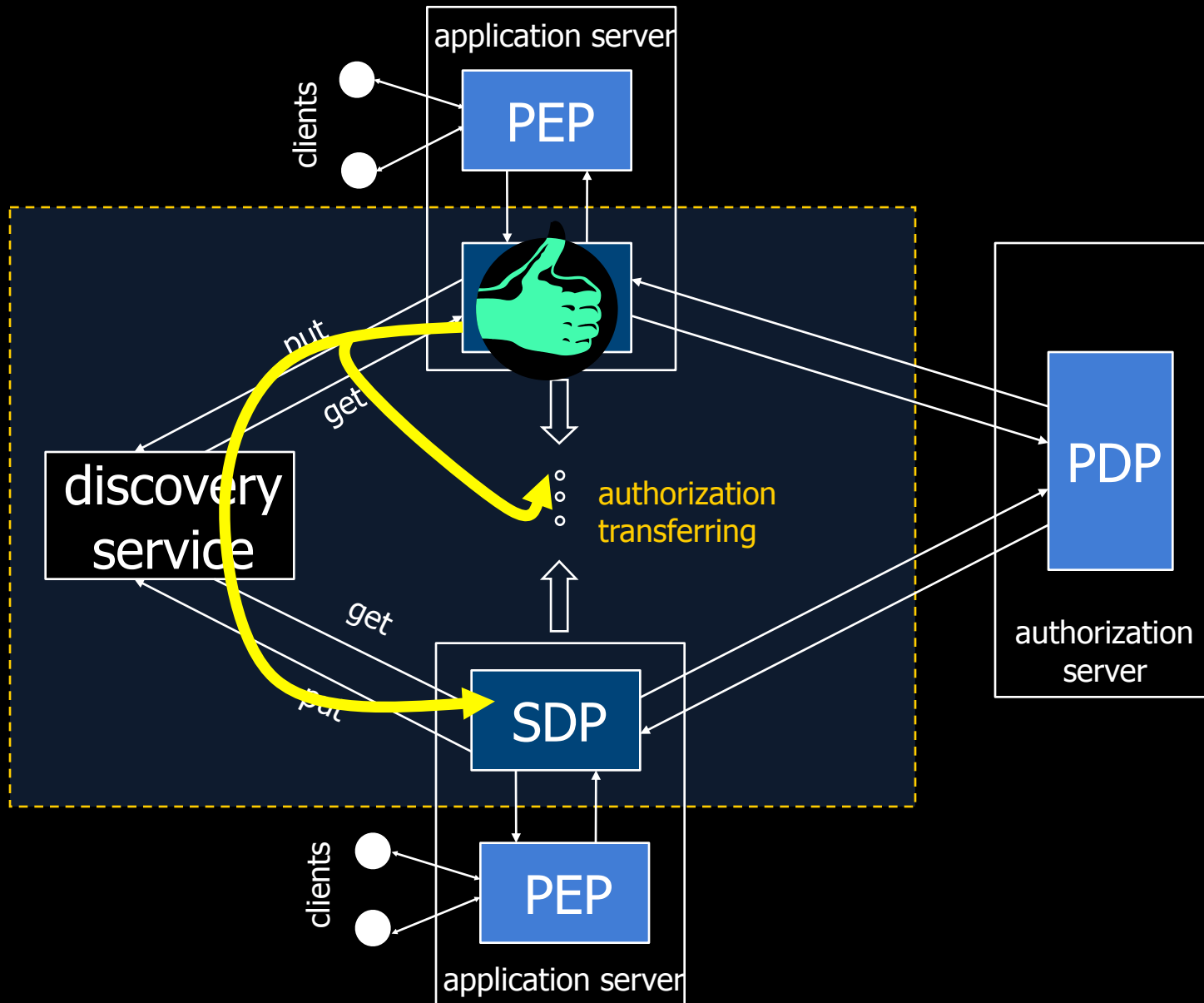


WAN



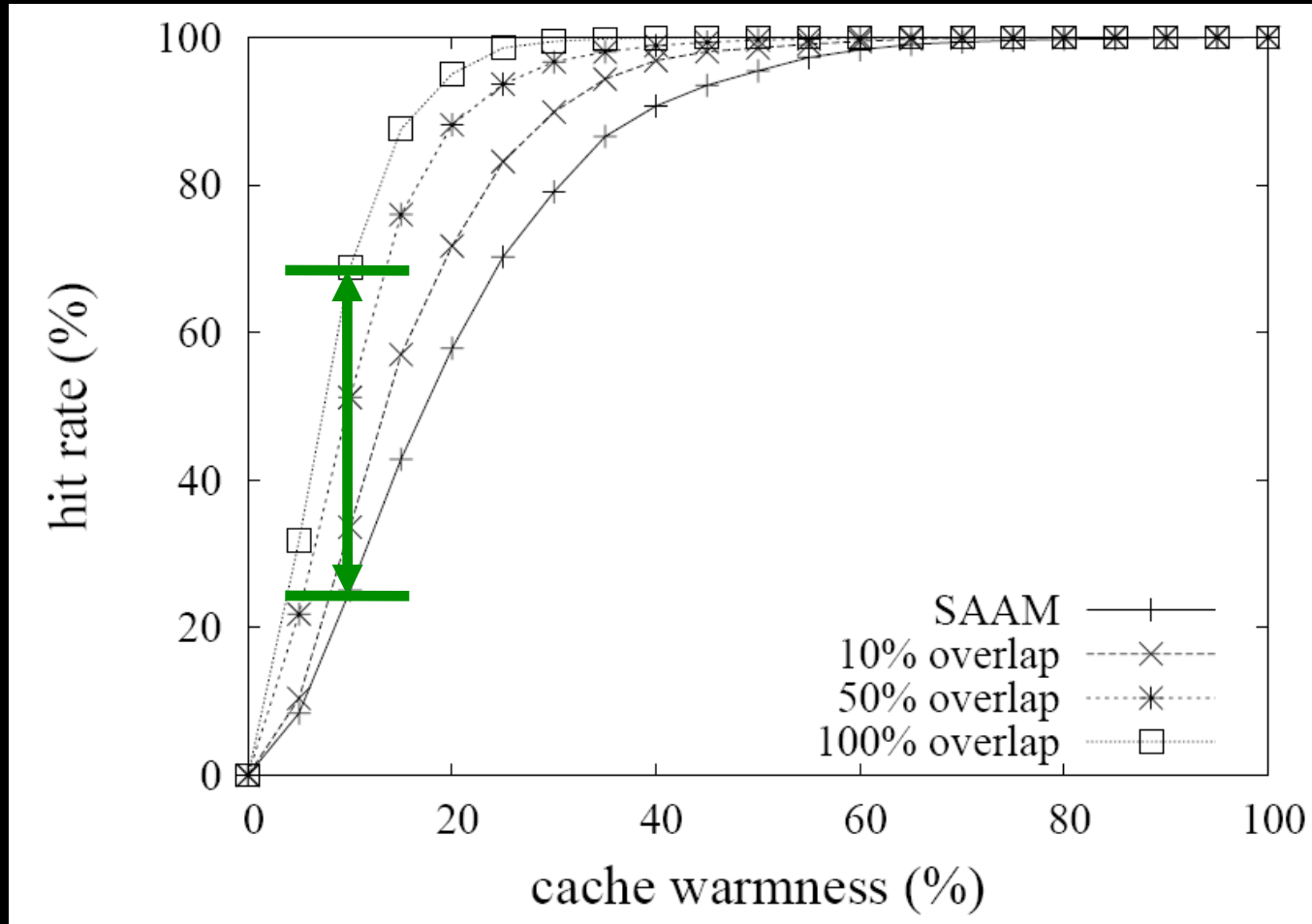
response time decreases with the number of requests
because more requests can be resolved by the local SDP

distributed and cooperative SAAM



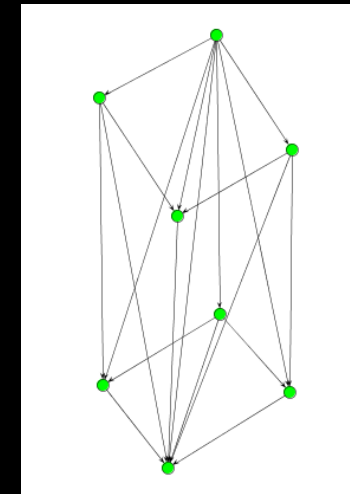
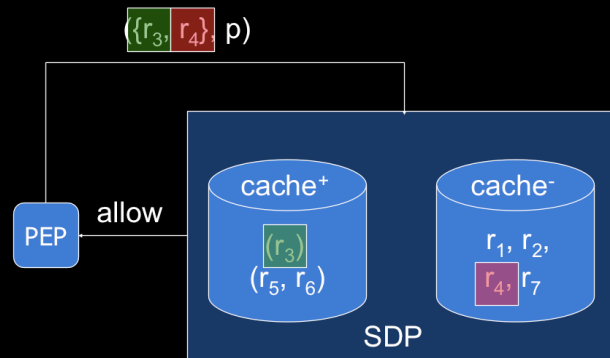
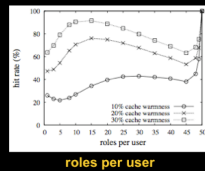
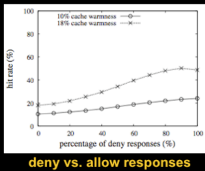
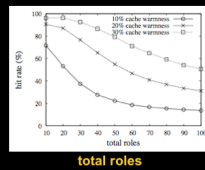
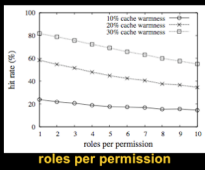
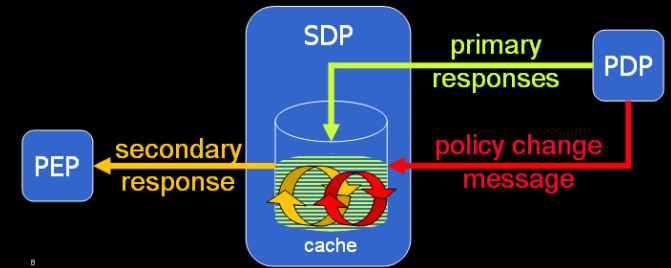
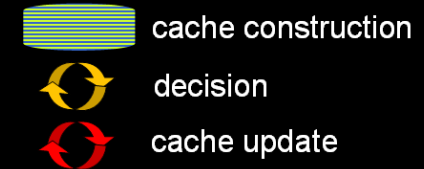
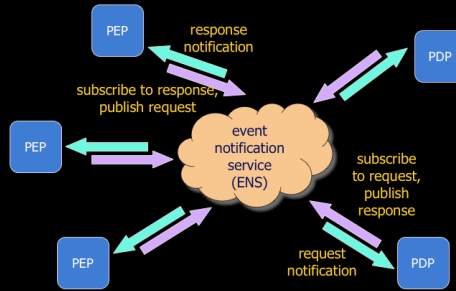
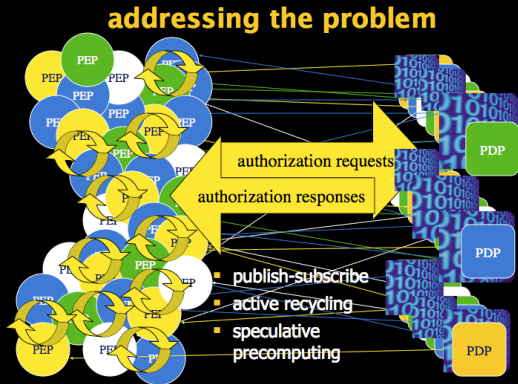
hit rate for distributed SAAM_{BLP}

5 SDPs' cooperation, uniform requests



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summary



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